

Introducing PLASSE Laboratory & its Current Research



Hanyang University, Ansan

Kyung-Goo Doh



Programming Languages **Application to** Software **Security &** Engineering

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Current Research Members

- professor : 1
- research professor : 1
- Ph.D. students : 4 (-3) (+1)
- M.S. students : 7 (-2) (+?)
- on-campus collaborators
- domestic collaborators
- international collaborators





Research Directions

- Theoretical research
 - develop theory and methodology
- Industrial applications
 - ✓ implement analysis engines and tools
 - transfer technology to software industry



Research Theme

- Semantics engineering
 - simple, user-friendly semantics meta-language
- String analysis
 - syntax and semantics analysis for dynamically generated strings
- Software security
 - static/dynamic detection of security vulnerabilities from source code
- Software maintenance
 - extraction of software properties and metrics from source code



Abstract Parsing

Joint work with Hyunha Kim & David A. Schmidt



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"Classic" String Analysis

- on statically generated string
- (scanning +) parsing





Parsing



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Dynamically Generated String



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Example: Database Application

public void printAddresses(String id) throws SQLException {
Connection con = DriverManager.getConnection("students.db");
String q = "SELECT * FROM address";
if (id != 0) q = q + "WHERE studentid=" + id;
ResultSet rs = con.createStatement().executeQuery(q);
while (rs.next()) { System.out.println.getString("addr")); }

taken from Christensen/Moeller/Schwartzbach's SAS2003 paper "Precise analysis of string expression" with minor modification

- This Java program parses and compiles OK.
- You can check if the dynamically generated SQL query parses OK at run-time.
- <u>Question</u>: Can we statically check if all the SQL queries generated by this program parse OK?



Previous Approach

• Analysis-then-parse





Previous Approach

• Analysis-then-regular_approximation-then-parse





Example



 Is the string generated at hot spot above conformed to the following reference grammar?

 $S \rightarrow "a" \mid "["S"]"$



Example: Previous String Analyzer



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Abstract Parsing

- Simultaneous analysis-and-parsing
 - statically analyze a program that dynamically generates strings, and, at the same time, parse the generated strings with the LR(k) reference grammar
- Abstract parse stacks as abstract string values
 - encode a generated string's context-free structure



Architecture of Abstract Parsing



Our abstract parser solves data-flow equations in the domain of abstract stack.

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LR(0) Parsing

Goto Controller for parser built from LR(0)-items for the reference grammar, $S \rightarrow "a" \mid "[" S "]"$



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"["



 $= S_5$



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 $= s_1 \cdot goto(s_1, "a")$

 $= s_1 \cdot qoto(s_1, S)$

 $= S_1 \cdot S_3$

 $= s_1 \cdot s_2$ (reduce: $S \rightarrow "a"$)





$$\begin{aligned} \mathsf{X}_0(\mathsf{s}_0) &= \operatorname{goto}(\mathsf{s}_0, \ "a") \\ &= \mathsf{s}_2 \quad (\operatorname{reduce:} \mathsf{S} \to "a") \\ &= \operatorname{goto}(\mathsf{s}_0, \mathsf{S}) \\ &= \mathsf{s}_5 \end{aligned}$$

$$\begin{aligned} X_1(s_0) &= ("[" \cdot "a")(s_0) \\ &= goto(s_0, "[") * "a" \\ &= s_1 * "a" \\ &= s_1 \cdot goto(s_1, "a") \\ &= s_1 \cdot s_2 \quad (reduce: S \to "a") \\ &= s_1 \cdot goto(s_1, S) \\ &= s_1 \cdot s_3 \end{aligned}$$

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 $= S_0 \cdot S_1$

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 $= S_0 \cdot S_1 \cdot S_1$

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Implementation





Conclusion

• Impact

- improve the precision on the syntax analysis of dynamically generated strings
- enhance the entire group of work based on previous stringanalysis technique.
- Future work
 - ✓ Abstract semantic processing on parsed strings.
 - type checking on dynamically generated strings
 - static analysis on dynamically generated strings



Discussion



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