

Class Analysis Workshop

Programming Research Lab. / SNU / Korea

January 17, 2009



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- Introduction
- Historical Overview
- Algorithm Overview
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 - k-CFA / Object-Sensitive
 - EPA / Refinement-Based
- Conclusion

Introduction

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Class Analysis Workshop
January 17, 2009



What is it?

to extract RTTI
from source code
statically

RTTI : runtime type information

What is it?

example

```
Parent * a;  
  
if(...)  
    a = new Child1();  
else  
    a = new Child2();  
  
a->do();
```

What is it?

example

```
Parent * a;  
  
if(...)  
    a = new Child1();  
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    a = new Child2();  
  
a->do();
```

What is it?

example

```
Parent * a;
```

```
if(...)
```

```
    a = new Child1();
```

```
else
```

```
    a = new Child2();
```

```
a->do();
```



What is it?

example

```
Parent * a;  
  
if(...)  
    a = new Child1();  
else  
    a = new Child2();  
  
a->do();
```

what is runtime type of a?

What is it?

example

```
Parent * a;  
  
if(...)  
    a = new Child1();  
else  
    a = new Child2();  
  
a->do();
```

what is runtime type of a?

Answer = {Child1, Child2}

What is it?

Similar things

Class Analysis : extract RTTI

Alias Analysis : extract aliased variables

Point-to Analysis : extract referencing location

Where to use?

Virtual Function Resolution (Compiler)

Call Graph Construction (Analysis)

Other Analysis (Analysis)

How to evaluate?

What is measure?

- size of points-to set
- # of call-edge
- # of resolved virtual call

-
- client analysis

Historical Overview

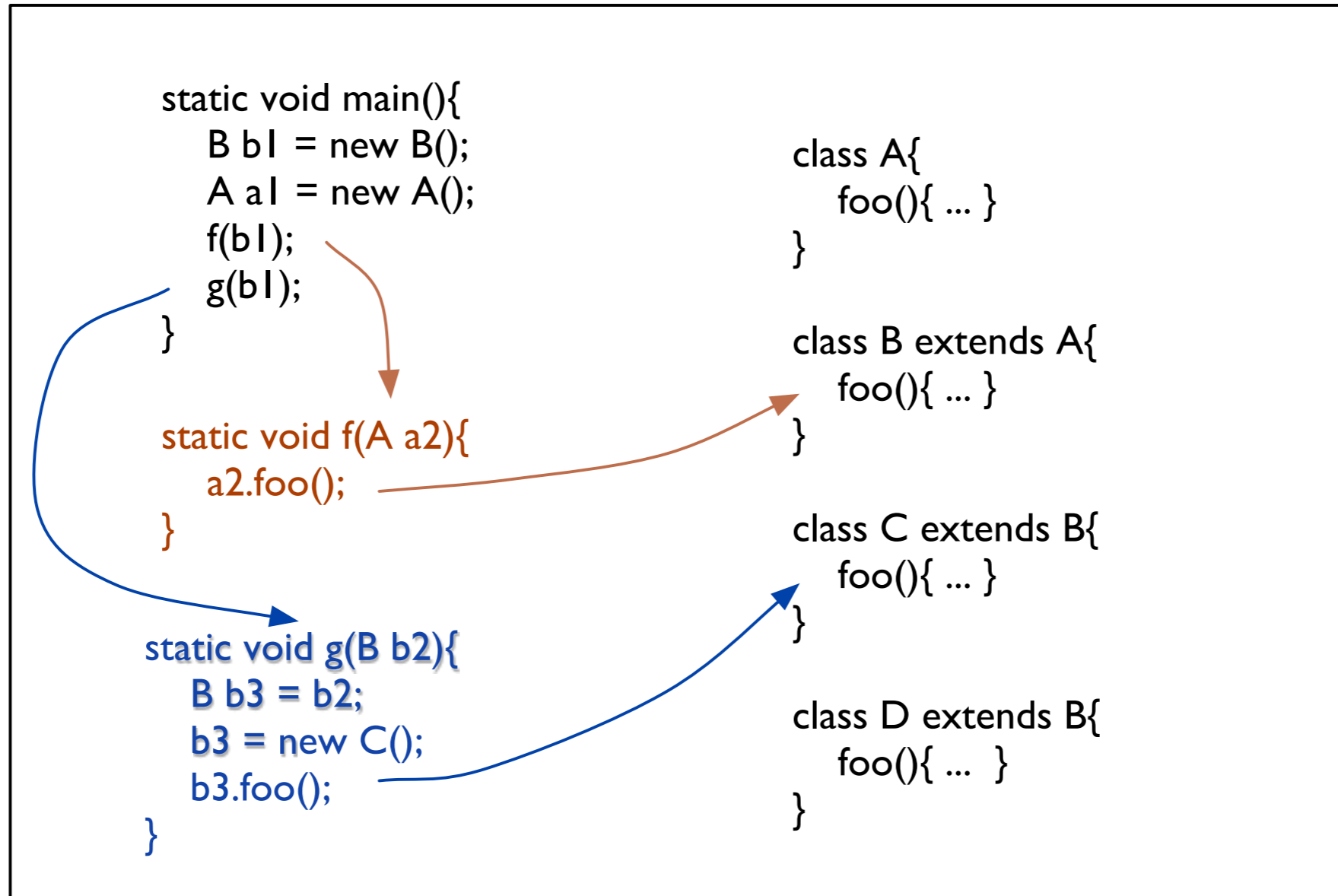
Part I

context-insensitive family

CHA / RTA / XTA / Anderson / 0-CFA

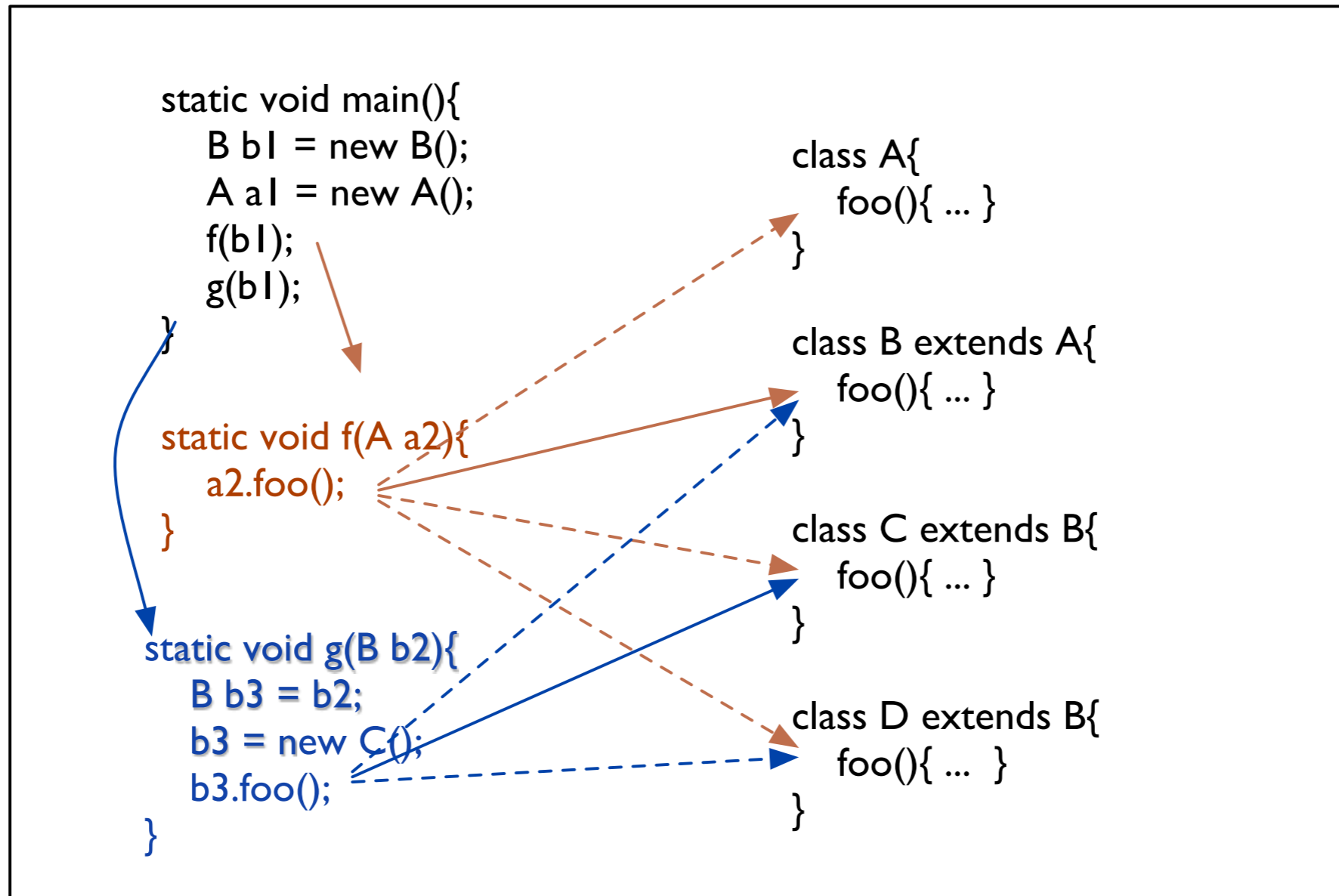
Example #1

cf Barbara G. Ryder, ETAPS 03



CHA

cf Barbara G. Ryder, ETAPS 03



J. Dean, D. Grove, C. Chambers,
“Optimization of OO Programs using Class Hierarchy”, ECOOP 95

CHA

Class Hierarchy Analysis

- calculate call graph using static type
- seems stupid
- resolve average 10-50% of Virtual Call
- recommend : working with other analysis
- $O(n)$

J. Dean, D. Grove, C. Chambers,
“Optimization of OO Programs using Class Hierarchy”, ECOOP 95

CHA

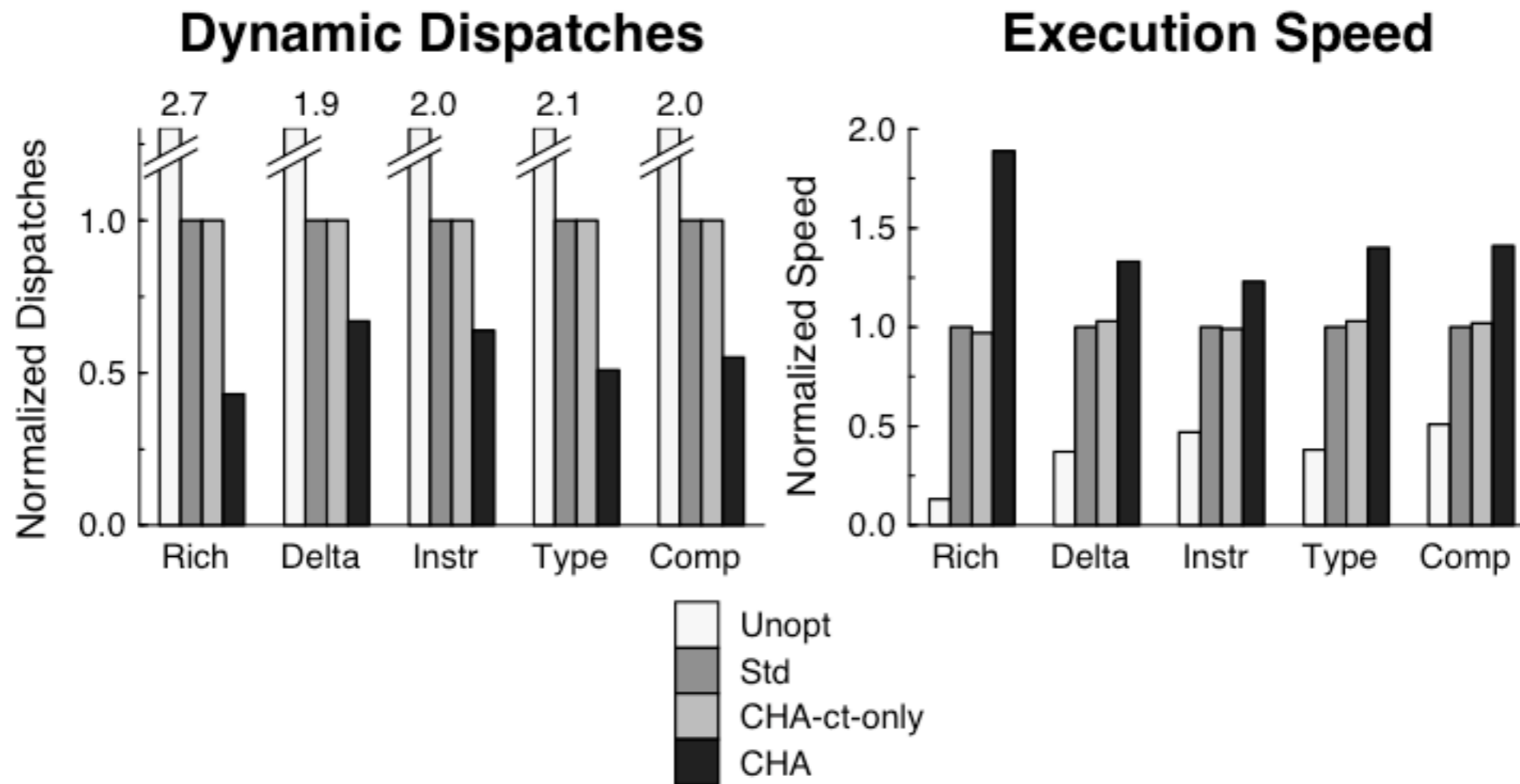
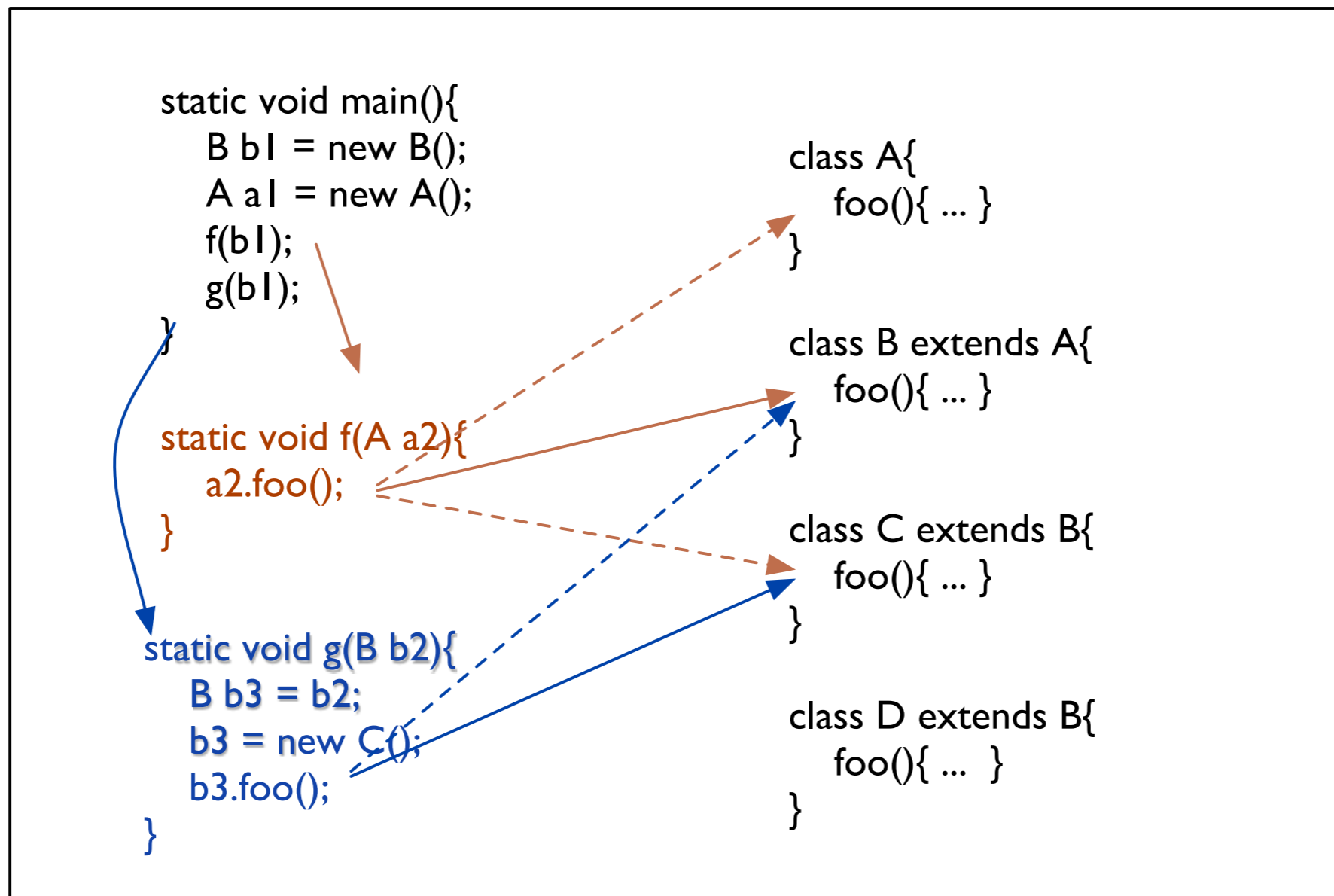


Figure 1: Number of dynamic dispatches and execution speed

chart from [Dean'95]

RTA

cf Barbara G. Ryder, ETAPS 03



D. Bacon and P. Sweeny,
“Fast Static Analysis of C++ Virtual Function Calls”, OOPSLA’96

RTA

Rapid Type Analysis

- calculate instantiated class set :“I”
- calculate call graph using “I” and CHA
- resolve 10-90% of virtual call
- $O(n)$

D. Bacon and P. Sweeny,
“Fast Static Analysis of C++ Virtual Function Calls”, OOPSLA’96



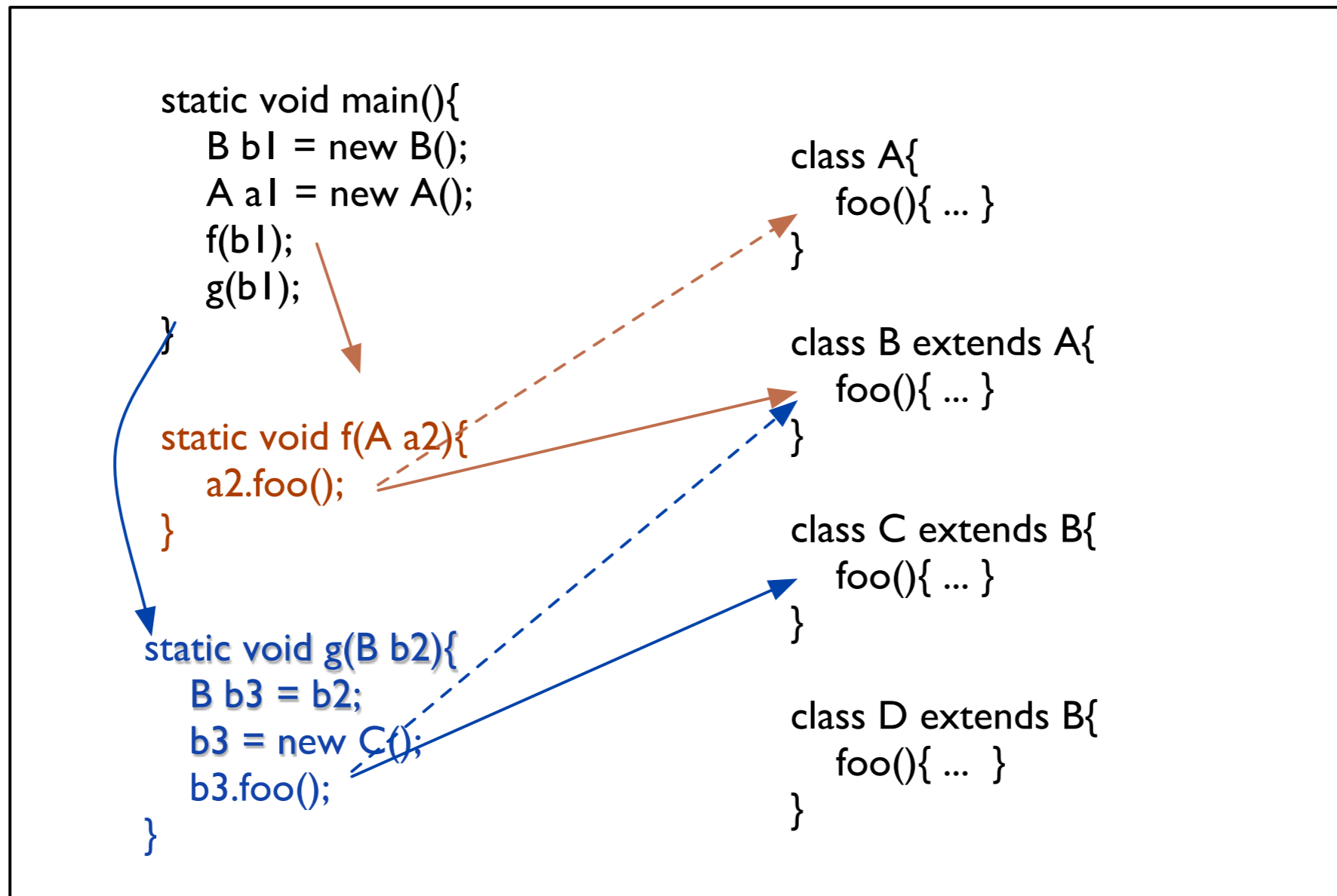
RTA

		Whole Application					
		Nodes		Edges		Callsites	
		Removed	(%tot.)	Removed	(%tot.)	Resolved	(%poly) (%tot.)
raytrace	pRTA	808	(46%)	3585	(39%)	292	(77%) (4.2%)
	oRTA	884	(51%)	4128	(45%)	300	(79%) (4.3%)
	DTA	925	(53%)	4375	(47%)	304	(80%) (4.4%)
	VTA	1026	(59%)	5200	(56%)	342	(90%) (4.9%)
	oRTA+VTA	1031	(59%)	5242	(57%)	342	(90%) (4.9%)
	VTA+VTA	1026	(59%)	5200	(56%)	342	(90%) (4.9%)
compress	pRTA	814	(51%)	3664	(45%)	293	(79%) (5.0%)
	oRTA	890	(56%)	4207	(52%)	301	(81%) (5.2%)
	DTA	926	(58%)	4418	(55%)	303	(82%) (5.2%)
	VTA	1033	(65%)	5214	(65%)	344	(93%) (5.9%)
	oRTA+VTA	1039	(65%)	5256	(65%)	346	(93%) (5.9%)
	VTA+VTA	1033	(65%)	5214	(65%)	344	(93%) (5.9%)
jack	pRTA	820	(44%)	3763	(34%)	313	(40%) (3.9%)
	oRTA	896	(48%)	4306	(39%)	321	(41%) (4.0%)
	DTA	924	(50%)	4475	(41%)	323	(41%) (4.1%)
	VTA	1027	(55%)	5719	(52%)	734	(94%) (9.2%)
	oRTA+VTA	1033	(55%)	5769	(53%)	735	(94%) (9.2%)
	VTA+VTA	1027	(55%)	5719	(52%)	734	(94%) (9.2%)

table from [Sundaresan'00]

XTA

cf Barbara G. Ryder, ETAPS 03



Tip and Palsberg, "Scalable Propagation-based Call Graph Construction Algorithms", OOPSLA'00

XTA

- natural extension of RTA
- one class set per method
- one class set per field
- propagate set through function call and return
- small improvement
- $O(n^3)$

Tip and Palsberg, “Scalable Propagation-based Call Graph Construction Algorithms”, OOPSLA’00

XTA

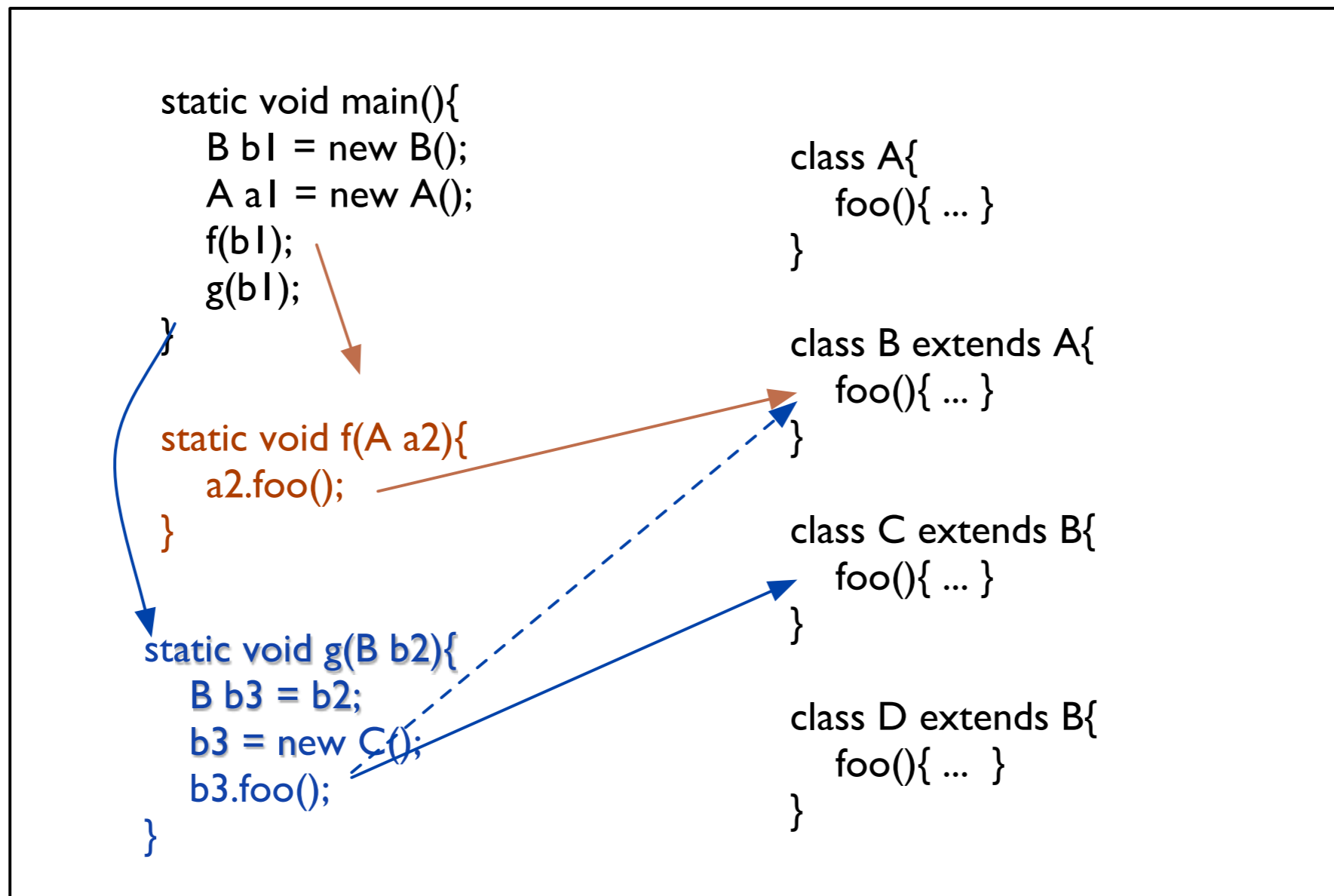
benchmark	RTA	MTA	FTA	XTA	(RTA-MTA)/RTA	(RTA-FTA)/RTA	(RTA-XTA)/RTA
Hanoi	400	386	382	379	3.5%	4.5%	5.3%
Ice Browser	1,594	1,594	1,593	1,588	0.0%	0.1%	0.4%
mBird	8,061	8,036	7,772	7,760	0.3%	3.6%	3.7%
Cindy	15,457	14,729	11,331	10,967	4.7%	26.7%	29.0%
CindyApplet	5,223	4,990	4,399	4,347	4.5%	15.8%	16.8%
eSuite Sheet	7,171	7,149	7,093	7,071	0.3%	1.1%	1.4%
eSuite Chart	14,669	14,648	13,857	13,771	0.1%	5.5%	6.1%
javaFig 1.43	5,128	5,064	4,963	4,961	1.2%	3.2%	3.3%
BLOAT	19,384	18,772	16,704	16,672	3.2%	13.8%	14.0%
JAX 6.3	7,053	7,018	6,904	6,895	0.5%	2.1%	2.2%
javac	13,154	13,154	13,115	13,113	0.0%	0.3%	0.3%
Res. System	46,130	45,944	44,792	44,412	0.4%	2.9%	3.7%
AVERAGE					1.6%	6.6%	7.2%

Number of edges in the call graph(RTA/XTA)

table from [Tip'00]

Anderson style

cf Barbara G. Ryder, ETAPS 03



L.Andersen, "Program Analysis and Specialization of the C Programming Language", PhD. thesis

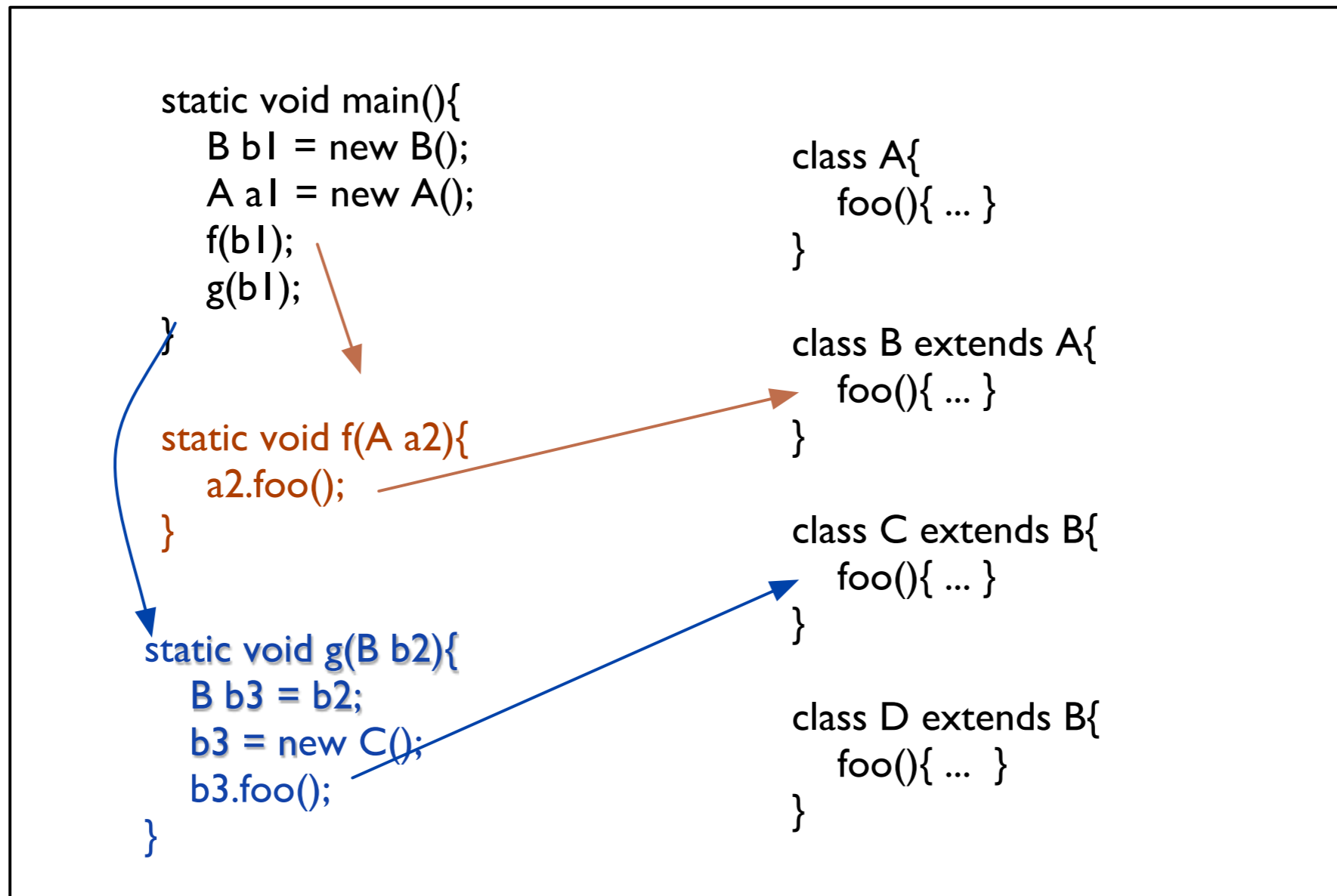
Rountev, A Milanova, B. Ryder, "Points-to Analysis for Java Using Annotated Constraints", OOPSLA'01

Anderson Style

- natural extension of XTA
- one class set per variable
- one class set per field
- distinguish object by allocation site
- $O(n^3)$

0-CFA

cf Barbara G. Ryder, ETAPS 03



Shivers, PLDI'88
Chambers, POPL'98

0-CFA

- flow-sensitive
- most precise in context-insensitive family
- (variable#) class set per program point
- $O(n^3)$
- expensive!

What to use?

- [Tip00]
 - Virtual Call Resolution : RTA
- [Sundaresan00][Lhotak06]
 - Call Graph Construction : Andersen Style

Historical Overview

Part 2

context-sensitive family

k-Context / Object-Sensitive

Example #2

cf Barbara G. Ryder, ETAPS 03

```
static void main(){
  D d1 = new D();
  if (...)
    (d1.f(new B())).g();
  else
    (d1.f(new C())).g();
}

public class D{
  public A f (A a1){ return a1; }
}

public class A{
  X xx;
  A (X xa){ this.xx = xa;}
}

public class B{
  B(X xb){C3 : super(xb);}
  public X f() { return this.xx;}
  static void main(){
    X x1, x2;
    B b1 = new B(new Y());
    B b2 = new B(new Z());
    x1 = b1.f();
    x1.g();
    x2.b2.f();
    x2.g();
  }
}
```

context-insensitive family cannot handle

k-call-site-sensitive

cf Barbara G. Ryder, ETAPS 03

```
static void main(){
  D d1 = new D();
  if (...)
    (d1.f(new B())).g();
  else
    (d1.f(new C())).g();
}

public class D{
  public A f (A a1){ return a1; }
}
```

```
public class A{
  X xx;
  A (X xa){ this.xx = xa;}
}

public class B{
  B(X xb){C3 : super(xb);}
  public X f() { return this.xx;}
  static void main(){
    X x1, x2;
    B b1 = new B(new Y());
    B b2 = new B(new Z());
    x1 = b1.f();
    x1.g();
    x2.b2.f();
    x2.g();
  }
}
```

John Plevyak, Andrew A. Chien, “Precise Concrete Type Inference for Object-Oriented Languages”, OOPSLA’94

John Whaley, Monica S. Lam, “Cloning-Based Context-Sensitive Pointer Alias Analysis Using Binary Decision Diagrams”, PLDI’04



k-call-site-sensitive

- unlimited call-site-string
- collapse SCC to one big function node
- flow-insensitive analysis
- using BDD to control context blow up
- implement Datalog to BDD compiler

k-call-site-sensitive

Benchmark	insens.	object-sensitive				call site			ZCWL
		1	2	3	1H	1	2	1H	
compress	2597	8.4	9.9	11.3	12.1	2.4	3.9	4.9	3.3
db	2614	8.5	9.9	11.4	12.1	2.4	3.9	5.0	3.3
jack	2870	8.6	10.2	11.6	11.9	2.4	3.9	5.0	3.4
javac	3781	10.4	17.7	33.8	14.3	2.7	5.3	5.4	
jess	3217	8.9	10.6	12.0	13.9	2.6	4.2	5.0	3.9
mpegaudio	2794	8.1	9.4	10.8	11.5	2.4	3.8	4.8	3.3
mtrt	2739	8.3	9.7	11.1	11.8	2.5	4.0	4.9	3.4
soot-c	4838	7.1	13.7	18.4	9.8	2.6	4.2	4.8	
sablecc-j	5609	6.9	8.4	9.6	9.5	2.3	3.6	3.9	
polyglot	5617	7.9	9.4	10.8	10.2	2.4	3.7	4.7	3.3
antlr	3898	9.4	12.1	13.8	13.2	2.5	4.1	5.2	4.3
bloat	5238	10.2	44.6		12.9	2.8	4.9	5.2	6.7
chart	7070	10.0	17.4		18.2	2.7	4.8		
jython	4402	9.9	55.9		15.6	2.5	4.3	4.6	4.0
pmd	7220	7.6	14.6	17.0	11.0	2.4	4.2	4.2	
ps	3875	8.7	9.9	11.0	12.0	2.6	4.0	5.2	4.4

Note: columns after the second column show multiples of the context-insensitive number.

Table 3. Number of equivalence classes of abstract contexts

graph from [Lhotak06]

k-call-site-sensitive

- [Lhotak,06]
 - even less precise than 2-call-site-sensitive
 - because of collapsed SCC

Object-Sensitive

cf Barbara G. Ryder, ETAPS 03

```
static void main(){
  D d1 = new D();
  if (...)
    (d1.f(new B())).g();
  else
    (d1.f(new C())).g();
}

public class D{
  public A f (A a1){ return a1; }
}

public class A{
  X xx;
  A (X xa){ this.xx = xa;}
}

public class B{
  B(X xb){C3 : super(xb);}
  public X f() { return this.xx;}
  static void main(){
    X x1, x2;
    B b1 = new B(new Y());
    B b2 = new B(new Z());
    x1 = b1.f();
    x1.g();
    x2.b2.f();
    x2.g();
  }
}
```

Ana Milanova, Barbara G. Ryder, “Annotated Inclusion Constraints for Precise Flow Analysis”, ICSM’03

Ana Milanova, “Light Context-Sensitive Points-to Analysis for Java”, PASTE’07

Object-Sensitive

- use receiver argument as context
- kind of bounded CPA [Agesen'95]
- about 2 time slower than Andersen Style
- more precise than k-call-site-sensitive

Comparison

[Lothak'06]

1. Object-Sensitivity + Heap Cloning
2. Object-Sensitivity
3. Call-Site + Heap Cloning
4. Call-Site
5. k-call-site
6. Adersen Style (context-insensitive)

Observe

[Lothak'06]

Object-Sensitive + Heap Cloning is best

Anderson Style : not very bad

Unrealizable Path caused by bounded-
context spoils precision

Modern Movement

- *goal : defeat Unrealizable Path*
- Equivalent Context [Xu08]
- Refinement-Based [Sridhan06]

Algorithm Overview

Contents

- CHA/RTA/XTA (정가을)
- 0-CFA (고윤석)
- Refinement-Based (공순호)
- Equivalent-Context (오학주)

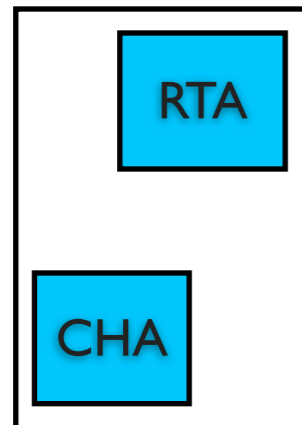
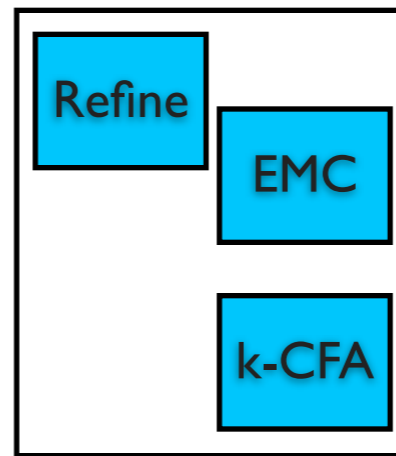
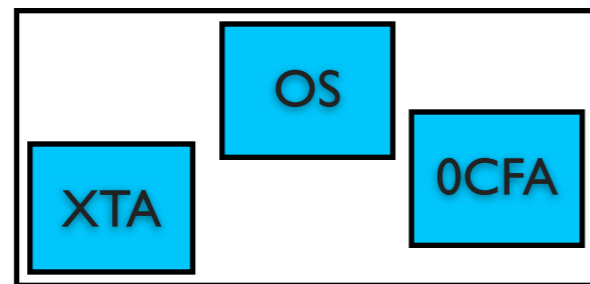
Conclusion

Categorize

precision



Naive



Implemented

Conceptual

time

What to use?

- Virtual Function Resolution :
RTA + Anderson Style
- Call Graph Construction : Object-Sensitive
- Pointer Characteristic Analysis :
Refinement-Based

Other Issues

- Demand-Driven?
- Unification-Based Analysis?

Thank you

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