



| Rosaec 2<sup>nd</sup> workshop

# Heap Data Structure Analysis using Abstract Parsing

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# Motivation



- ❖ **To determine the correctness of heap data structures generated during the program execution.**



# Approach



- ❖ Transform the procedure of generating data structures into generating strings.
- ❖ Then, parsing those strings to determine correctness of data structures.
- ❖ To achieve our goal, we'll use 'Abstract Parsing' technique.



# Reference



- ❖ **Kyung-Goo Doh, Hyunha Kim and David Schmidt. *Abstract parsing: static analysis of dynamically generated string output using LR-parsing technology*. SAS 2009: The 16th International Static Analysis Symposium, LA, August 7-9, 2009. (to appear)**

# Background



## ❖ What is the 'Abstract Parsing?'

```
❖ $s = "SELECT name WHERE class = ".$class."  
FROM students;";  
$result = db_query($s);
```

**To check syntactic correctness or  
harmlessness of \$s**

# Background (cont.)



## ❖ Abstract Parsing : Idea

- Instead of executing the program and parsing the result,

$$\mathcal{V}_0 e \Sigma = \{c_1, c_2, \dots, c_n\} \quad parse(c_i) = O/X$$

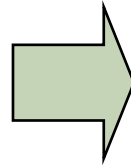
- Execute the program on abstract semantics using LR parsing technology.

$$\hat{\mathcal{V}}_0 e \hat{\Sigma} = \{O, X\}$$

# Background (cont.)



```
x = 'a'  
r = '']  
while ...  
    x = '[' . X . R  
print x
```



```
X0 = a  
R = ]  
X1 = X0 U X2  
X2 = [ . X1 . R  
X3 = X1
```

Each flow equations is interpreted as functions  $P \rightarrow P$   
(input parse stack  $\rightarrow$  output parse stack)



# Grammar



## ❖ Grammar of correct binary tree

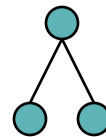
$$S \rightarrow \begin{array}{l} n \text{ ( ) ( )} \\ | n \text{ ( S ) ( S )} \end{array}$$

(\*type S1\*)

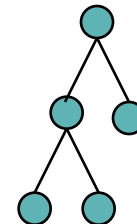
(\*type S2\*)

## ❖ Example

$n \text{ ( n ( ) ( ) ) ( n ( ) ( ) )}$



$n \text{ ( n ( n ( ) ( ) ) ( n ( ) ( ) ) ) ( n ( ) ( ) )}$



# Grammar (cont.)



## ❖ Grammar of incorrect binary tree

<b>I</b>	<b>→</b>	<b>n (S) ()</b>	<b>(*type I1*)</b>
		<b>  n () (S)</b>	<b>(*type I2*)</b>
		<b>  n () (I)</b>	<b>(*type I3*)</b>
		<b>  n (I) ()</b>	<b>(*type I4*)</b>
		<b>  n (S) (I)</b>	<b>(*type I5*)</b>
		<b>  n (I) (S)</b>	<b>(*type I6*)</b>
		<b>  n (I) (I)</b>	<b>(*type I7*)</b>

# Language



**c** → **x := e**  
| **x.1 := e**  
| **x.2 := e**  
| **c1 ; c2**  
| **if e c1 c2**  
| **while e c**

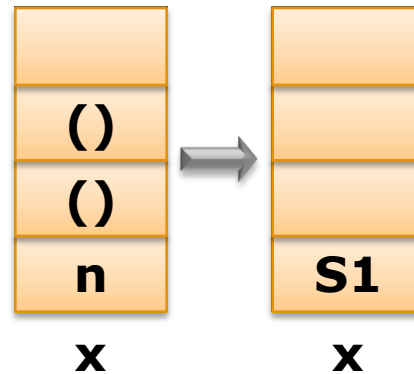
**e** → **x**  
| **malloc**  
| **nil**



# Example : Case 1

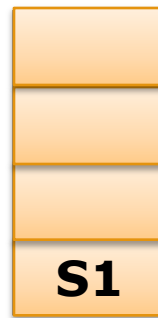


```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
}
```

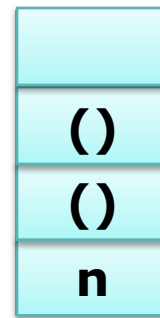


# Case 1 (cont.)

```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
}
```



x



y

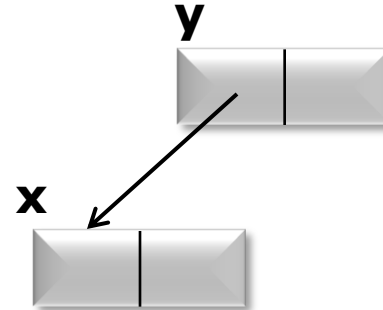


y

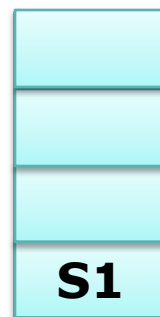
# Case 1 (cont.)



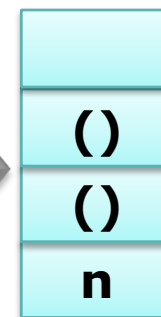
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x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
}
```



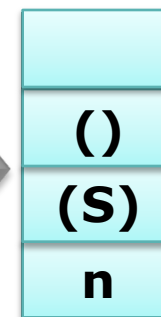
x



y



y



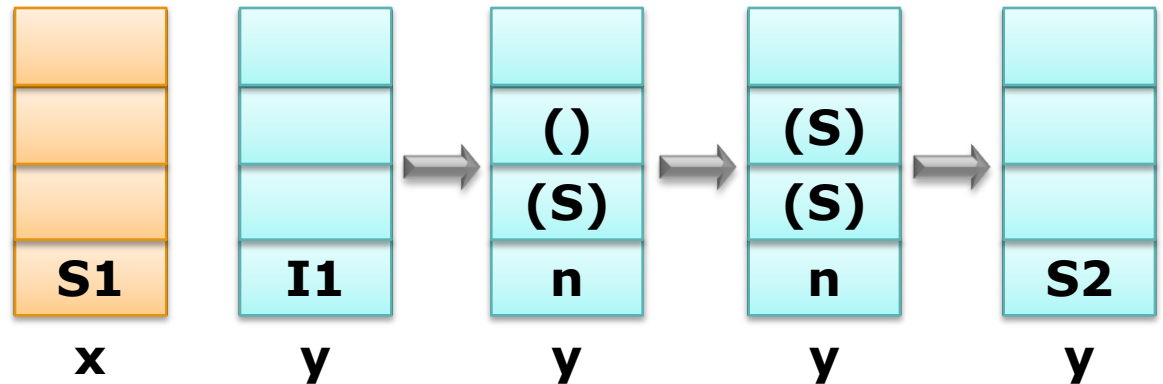
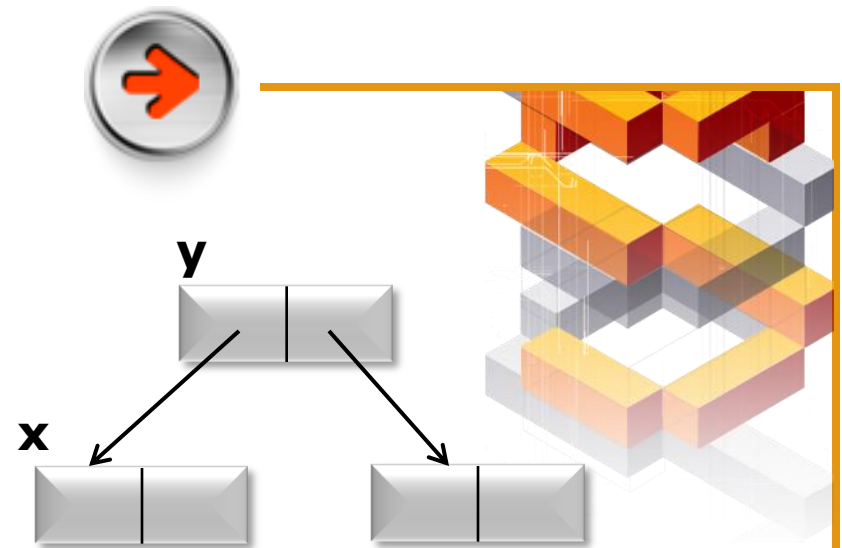
y



y

# Case 1 (cont.)

```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
}
```

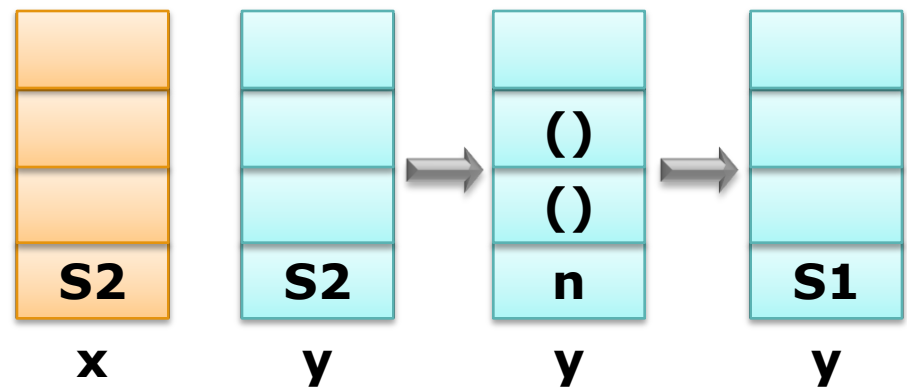
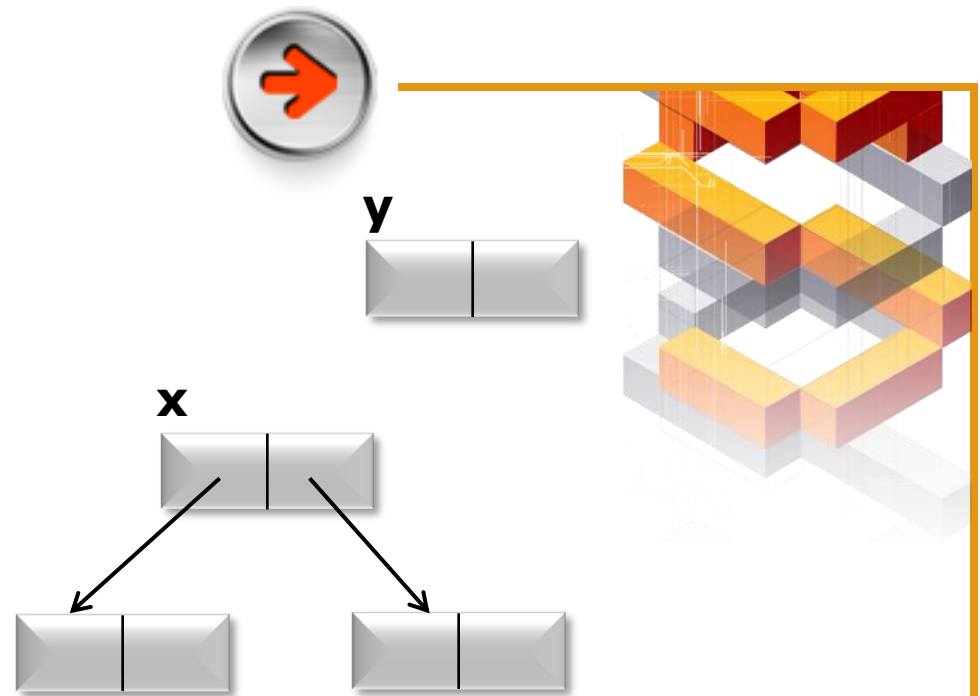






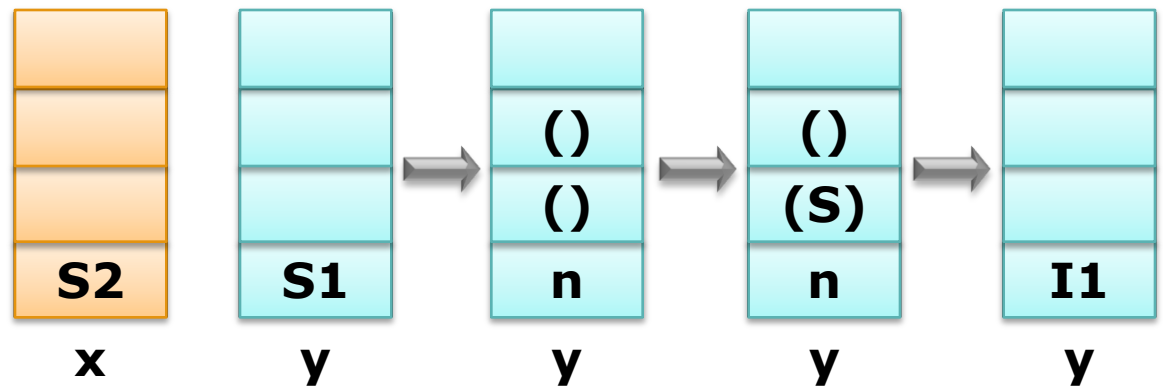
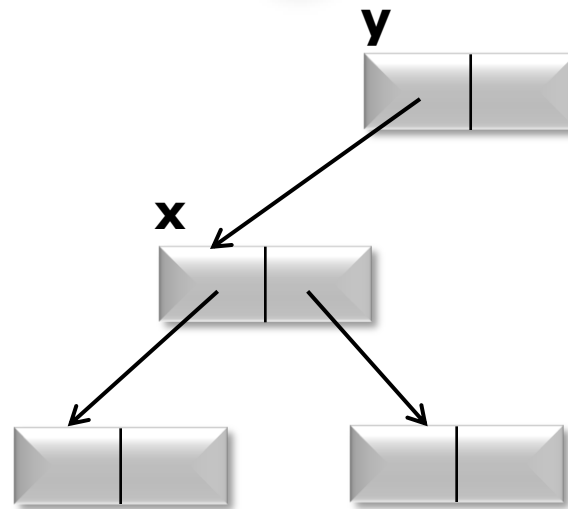
# Case 1 (cont.)

```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
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```



# Case 1 (cont.)

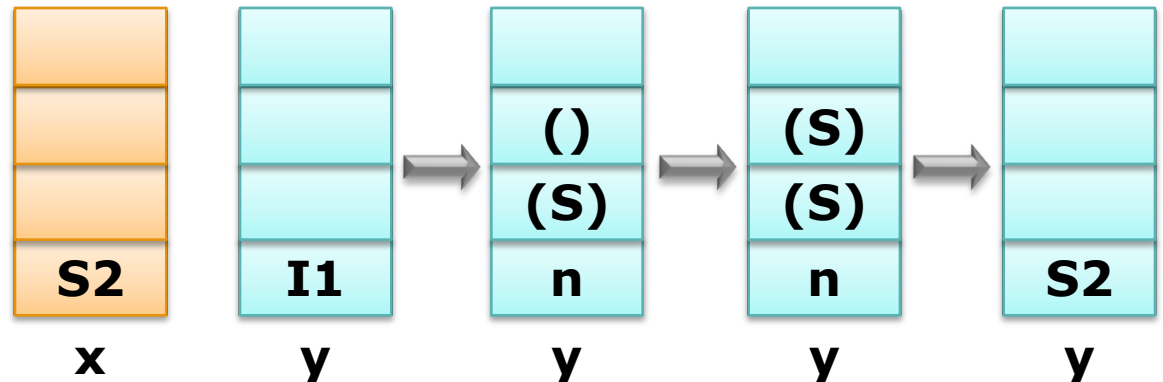
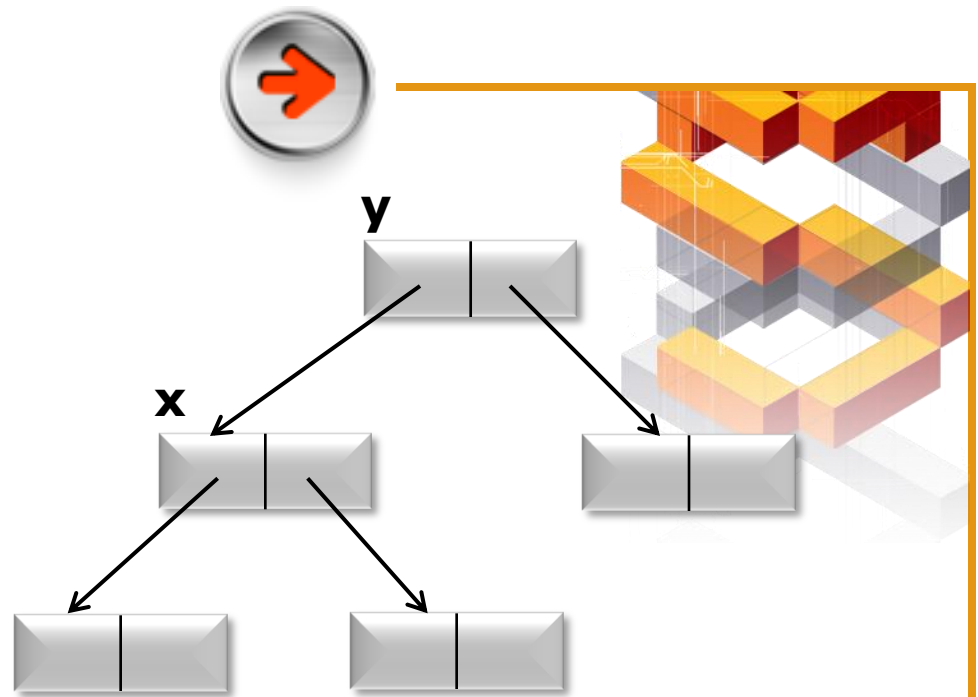
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x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
}
```



# Case 1 (cont.)

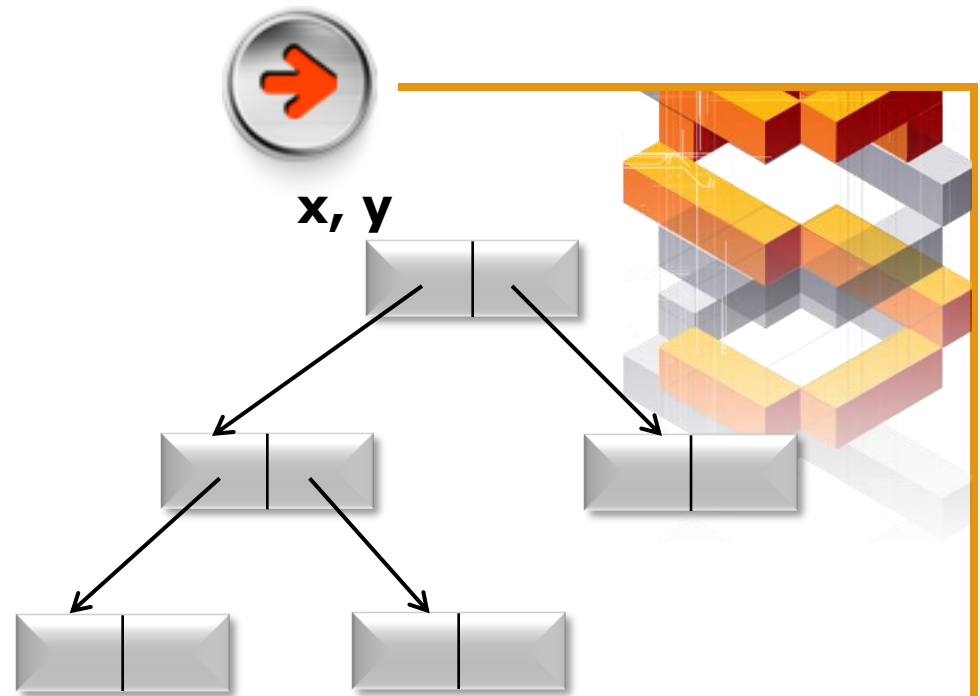
```

x := malloc;
while (x)
{
  y := malloc;
  y.1 := x;
  y.2 := malloc;
  x := y
}
  
```

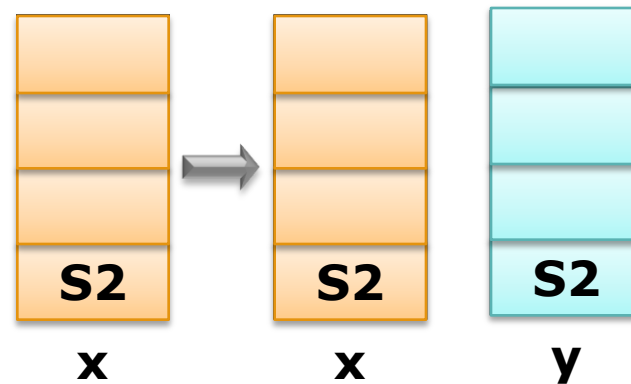


# Case 1 (cont.)

```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  y.2 := malloc;  
  x := y  
}
```



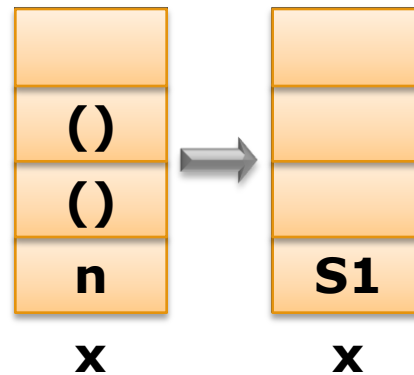
Reached fix point !!



# Case 2



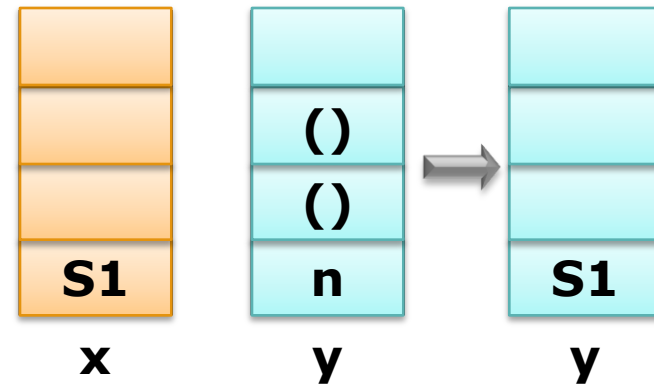
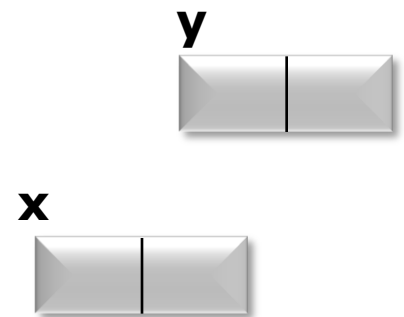
```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  x = y  
}
```



# Case 2 (cont.)



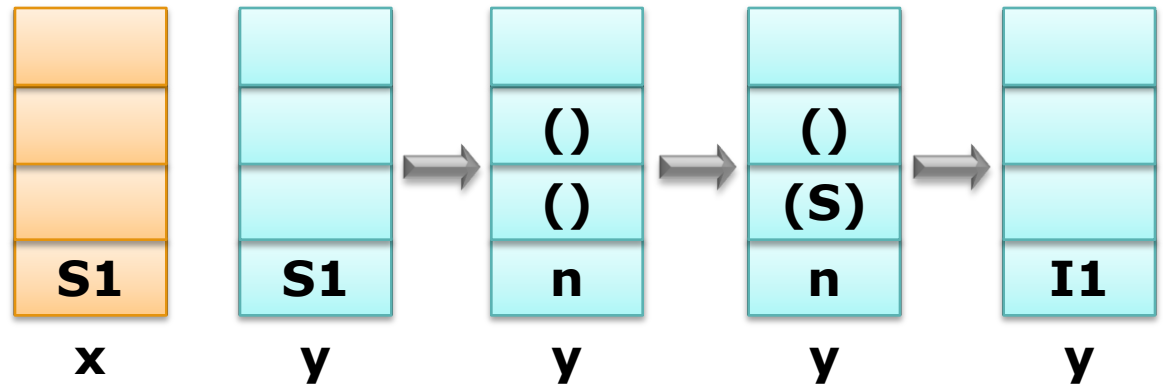
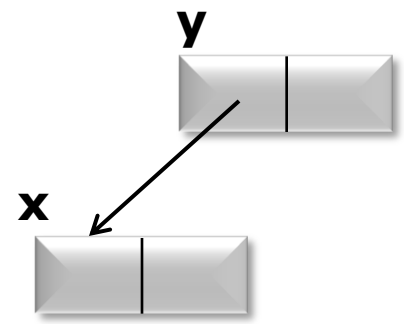
```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  x = y  
}
```



# Case 2 (cont.)

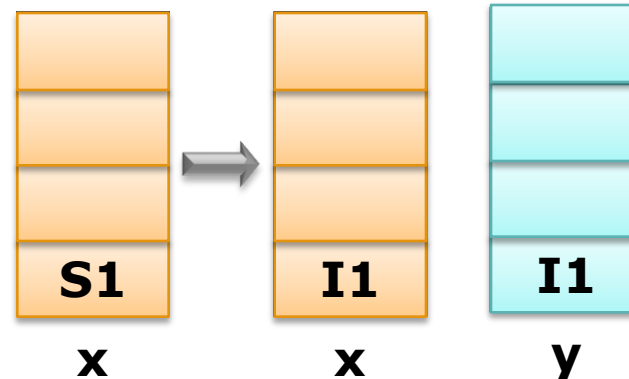
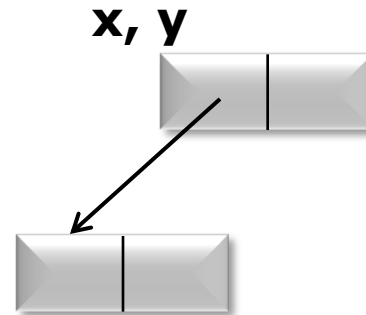


```
x := malloc;  
while (x)  
{  
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  y.1 := x;  
  x = y  
}
```



## Case 2 (cont.)

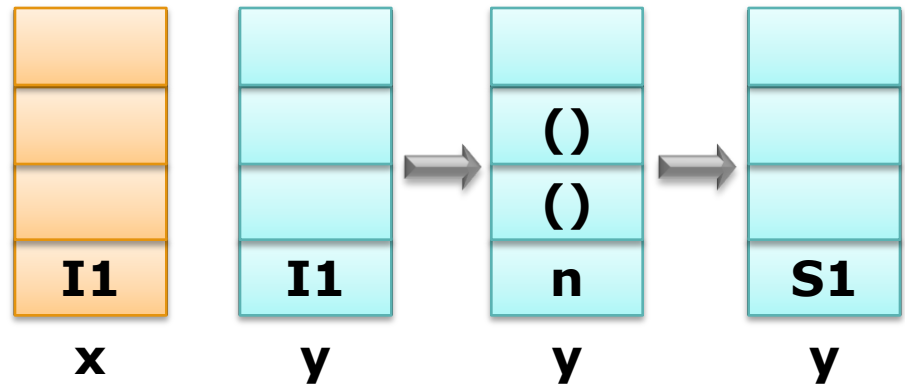
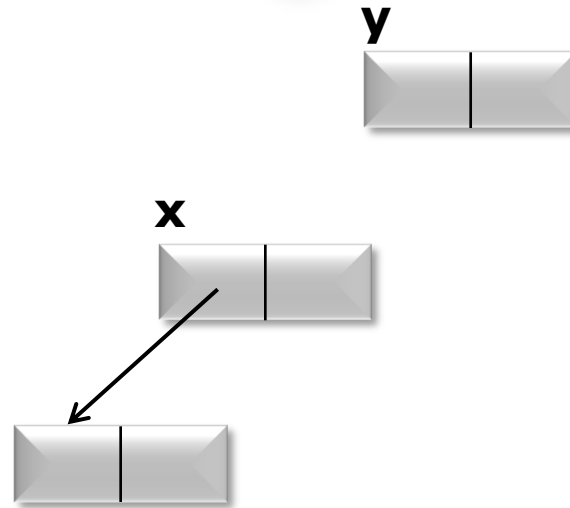
```
x := malloc;  
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{  
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  y.1 := x;  
  x = y  
}
```





# Case 2 (cont.)

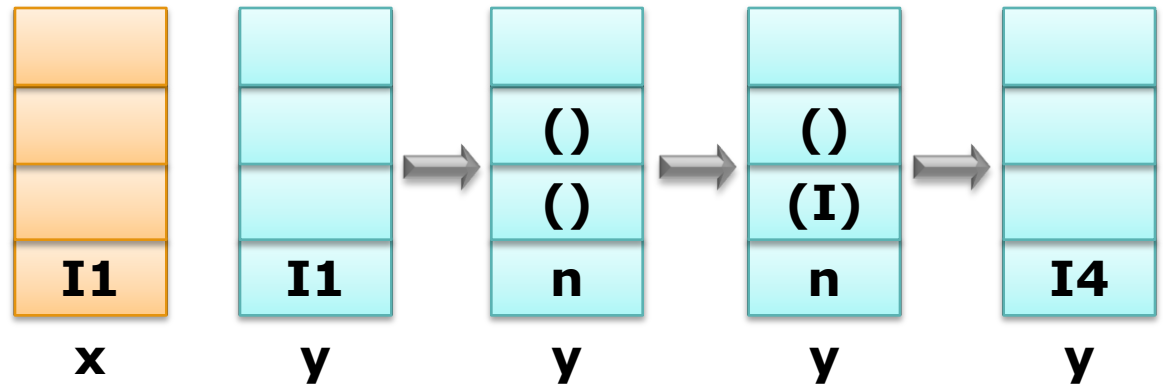
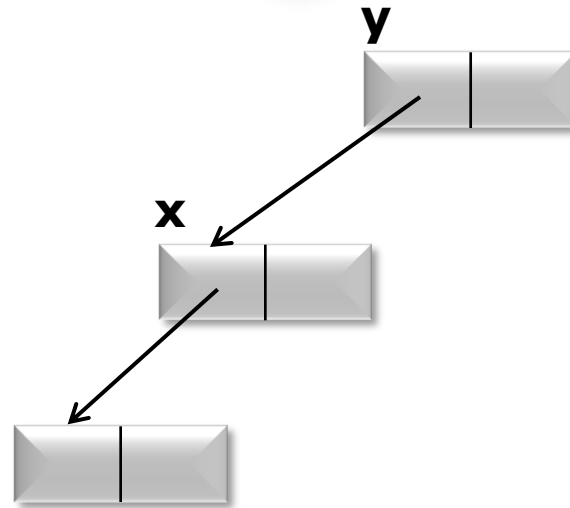
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x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  x = y  
}
```



# Case 2 (cont.)

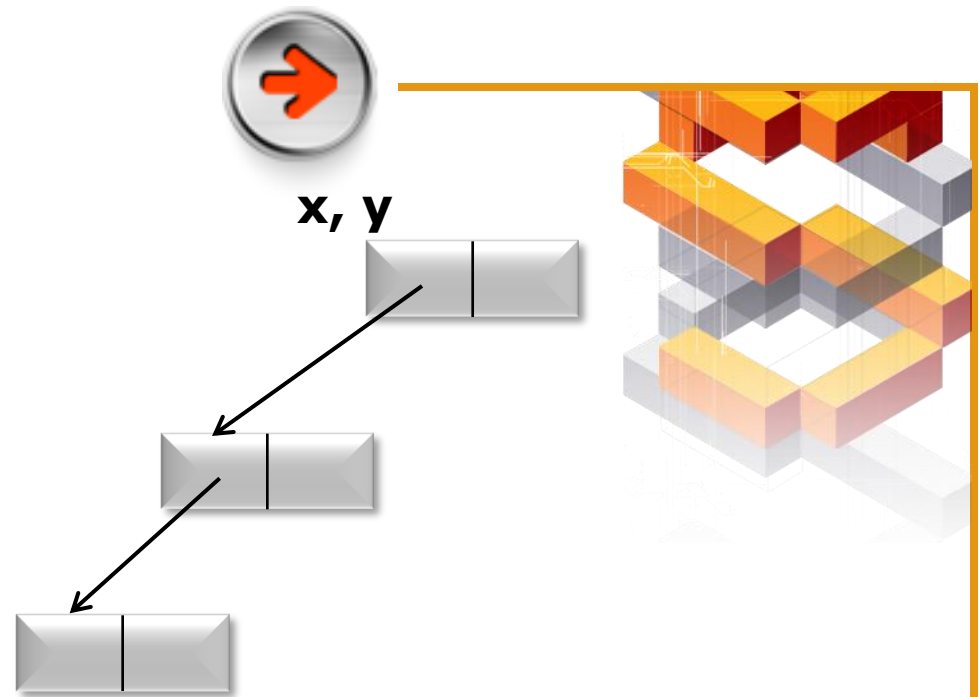


```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  x = y  
}
```

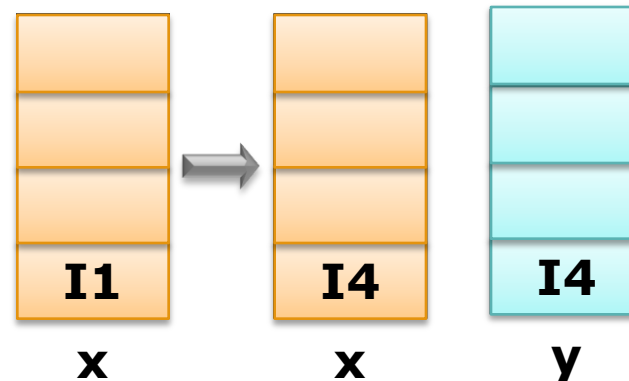


## Case 2 (cont.)

```
x := malloc;  
while (x)  
{  
  y := malloc;  
  y.1 := x;  
  x = y  
}
```



Reached fix point !!



# Future work



## ❖ Language expansion

### (Make followings possible)

- Direct memory access, sub-tree dereferencing, cyclic update.

## ❖ Grammar expansion

- Doubly linked list, linked list containing linked lists etc.



**Thank You !**

