



Introduction to KNU DBLab and Robust Tuple Extraction

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Recent Publications

Query Optimization

- Progressive Optimization [SIGMOD'07, VLDB'06]
- Parallelizing Query Optimization [SIGMOD'09, VLDB'08]

Query Processing

- Similarity Search in Large Time-Series Data [VLDB'07, SIGMOD'01]
- Applying DB Query Processing to Massive Amount of Streaming Data
 - Streaming XML Processing [VLDB'08, WWW'07]
 - Streaming Moving Object Processing [IEEE TKDE'09]



1 conference paper published (SIGMOD 2009)

I conference paper submitted

Tuple Extraction from Web Pages

 Extracting tuples from HTML pages has been an important issue in various applications

Applications

- -Web data integration
- -E-commerce market monitoring
- -Mashups

An Example Web Page



A Web Page at Time T₀



Motivation

The current state-of-the-art extraction systems are very vulnerable to small changes on the web

A robust extraction solution is needed!

Dependency-Aware Reordering for Parallelizing Query Optimization (SIGMOD 2009)

An example query

SQL: SELECT * FROM Professor P, Course C, Student S WHERE P.pid = C.pid AND S.sid = C.sid



Only need to specify what to find!

Query Execution Plans (QEPs)

Plan = Query Execution Plan (QEP) = Access Plan



Example: Query Optimization



DP Query Optimization

 Query optimization exploits dynamic programming (DP) to avoid generating redundant QEPs.

 Query optimization times using DP can increase significantly as the number of joins in a query increases.

Query Complexity vs. CPU Performance



Moore's Law outperforms query complexity

Multi-Core Wave



Motivation

- Parallelizing query optimization significantly delays the need to rely on suboptimal heuristics
- Our previous work on parallelizing query optimization [VLDB08] has some limitations
 - No support for recently developed enumeration algorithms
 - Static search space allocation can lead to very unbalanced workloads
 - Does not fully exploit parallelism due to local memo merge

Example of DPccp Optimizer

SELECT *	
FROM $R_1 q_1, R_2 q_2,$	
$R_3 \ q_3, R_4 \ q_4$	
WHERE $q_1.a_2 = q_2.a_1$ a	and
$q_1.a_3 = q_3.a_1$ a	and
$q_1.a_4 = q_4.a_1$ a	and
$q_2.a_3 = q_3.a_2$ a	and
$q_2.a_4 = q_4.a_2$ a	and
$q_{3.a_4} = q_{4.a_3}$	



1 (a)
$\begin{array}{c} 1. & (q_4) \\ 2. & (q_3) \end{array}$
$\overline{3}$. (q_2)
4. (q_1)
5. (q_3, q_4)
$\begin{array}{ccc} 0. & (q_2, q_4) \\ 7 & (q_2, q_2) \end{array}$
8. (q_2, q_3)
9. (q_4, q_2q_3)
10. (q_3, q_2q_4)
11. (q_1, q_4)
$12. (q_1, q_3)$ 13 (q_1, q_2, q_4)
$14. (q_1, q_3q_4)$
15. (q_1, q_2q_3)
16. (q_1, q_2q_4)
$17. (q_1, q_2q_3q_4)$
$18. (q_4, q_1q_2)$ 19 (q_2, q_4q_2)
$20. (a_1a_2, a_3a_4)$
$21. (q_4, q_1q_3)$
22. (q_2, q_1q_3)
23. (q_1q_3, q_2q_4)
24. $(q_4, q_1q_2q_3)$
$25. (q_3, q_1q_4)$ $26. (q_2, q_1q_4)$
$27. (q_1q_4, q_2q_3)$
28. $(q_3, q_1q_2q_4)$
29. $(q_2, q_1q_3q_4)$



e change the order?

Problem Statement

Develop a generic framework for parallelizing any bottom-up query optimizer

Overview of Solution

- A totally ordered sequence of pairs of quantifier sets is generated in a streaming fashion
- Buffer a fixed number of pairs and delay plan generation for those pairs
- Perform dependency-aware reordering
 - Convert the total order into a partial order over groups
 - Execute parallel group topological sort

 To maximize parallelism, we propose three novel optimizations

A Partial Order



Example



Experiment 1

To determine how much the three optimization techniques contribute to maximize the parallelism of DPEGeneric



Star query with 20 quantifiers

THANK YOU!

Any Question?