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APPLYING CONCOLIC
TESTING TO SOFTWARE
WRITTEN IN MULTI-
STAGED LANGUAGE

지난 이야기

악성 자바스크립트
코드 검출을 위한
분석기 고안

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Mal-JavaScript를 찾아서

기

2009년 파주에서 열린
2nd ROSAEC Workshop에서...

지난 이야기



악성 JavaScript 코드 (예)

```
delob5=new Array(7,84,89,81,94,88,79,27,82,95,6,67,86,87,79,90,73,92,94,79,27,88,87,90,72,72,82,95,6,25,120,119,104,114,/* ... */5);

function goMDAC(){
Try{
var PEELt6 = Qy29Nd.CreateObject(dddec("Shell.Application"), '');
JB7Ebp.open("GET", "http://2.gooanal.net/sis/getexe.php?h=11", false);
JB7Ebp.send();
LoWMFJ.Write(JB7Ebp.responseBody);
Frogxa = "\087akbY.exe";
```

```
eval(dddec("PEELt6.ShellExecute(Frogxa);"));
```

5

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지난 이야기



목표

Deobfuscate된 malicious JavaScript 코드에 대해
이 코드가
어떤 코드인지를
의미 기반 분석
수행 전에 판단하는 분석기

7

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어떤 코드가 동적으로 생성, 수행될지를 분석할 필요

Technical Memo

September 8, 2009 ROSAEC-2009-06 3

September 8, 2009 ROSAEC-2009-06

ROSAEC center
Research On Software Analysis for Error-free Computing
소프트웨어 무결성 연구센터 ROSAEC EFC

ROSAEC MEMO
2009-05
September 4, 2009

A Control Flow Analysis for 2-staged Programming Languages

Taeksu Kim¹, Chunwoo Lee¹, Kiljoo Lee¹, Soohyun Baik¹, and Kwangkeun Yi²
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A bstract

As a program written in multi-staged language can generate and execute code fragments in execution time, it is hard to predict what code fragments will be generated and run in execution time. Therefore we need a new method to predict what code fragments would be generated in execution time to analyze a control flow of software.

In this article, we present static analysis which detects code fragments generated in execution time using abstract interpretation. Moreover we prove the correctness of analyzer.

1 Introduction

Multi-staged language is a language which can generate and execute new program codes in execution time. The key aspect of multi-staged languages is to have code templates (program fragments) as first-class objects [1]. Recently multi-staged languages are used widely because they can make software development faster and prevent side effects occurred with single-staged language. Many Web programming languages(e.g. Ruby, Python) or script languages(e.g. PHP, Lisp, JavaScript) adopt multi-staged features in their specification so that they support rapid development.

Even if multi-staged languages are useful for rapid software development, it is hard to estimate what code fragments will be generated and executed. Therefore, the estimation of a control flow of software written in multi-staged language is hard, too. To analyze the control flow of the software, it is needed to analyze how many and what code fragments would be generated in execution time.

In this article, we design a static analysis for detecting code fragments generated in execution time using abstract interpretation[2] and prove the correctness of analyzer. We expect that this analyzer could help software developers to analyze a control flow of multi-staged programming language.

2 Language

Figure 1 shows a syntax of a simple 2-staged language. In the language, a program consists of expressions. Functions and code fragments are the first-class objects in the language. This language use quasi-quotations[3] to generate and execute code fragments in the execution time. Boxing(¹) operator makes a subexpression as a code fragments at one-stage and unboxing(·) operation substitutes a subexpression by a value evaluated at zero-stage. Run operation evaluates a code fragment. Boxing and unboxing operators are labeled by unique label and alphabet respectively. Detailed semantics are shown in the next chapter.

To decrease the complexity of the analysis we make some assumptions on the language as follows.

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A Control Flow Analysis for 2-staged Programming Languages

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September 2009.

JavaScript Strikes Back



문자열 합성 기반의 코드 생성

A New Hope

SE에서 기존에 논의되던 분석 방법을
Multi-staged 언어에 적용할 때

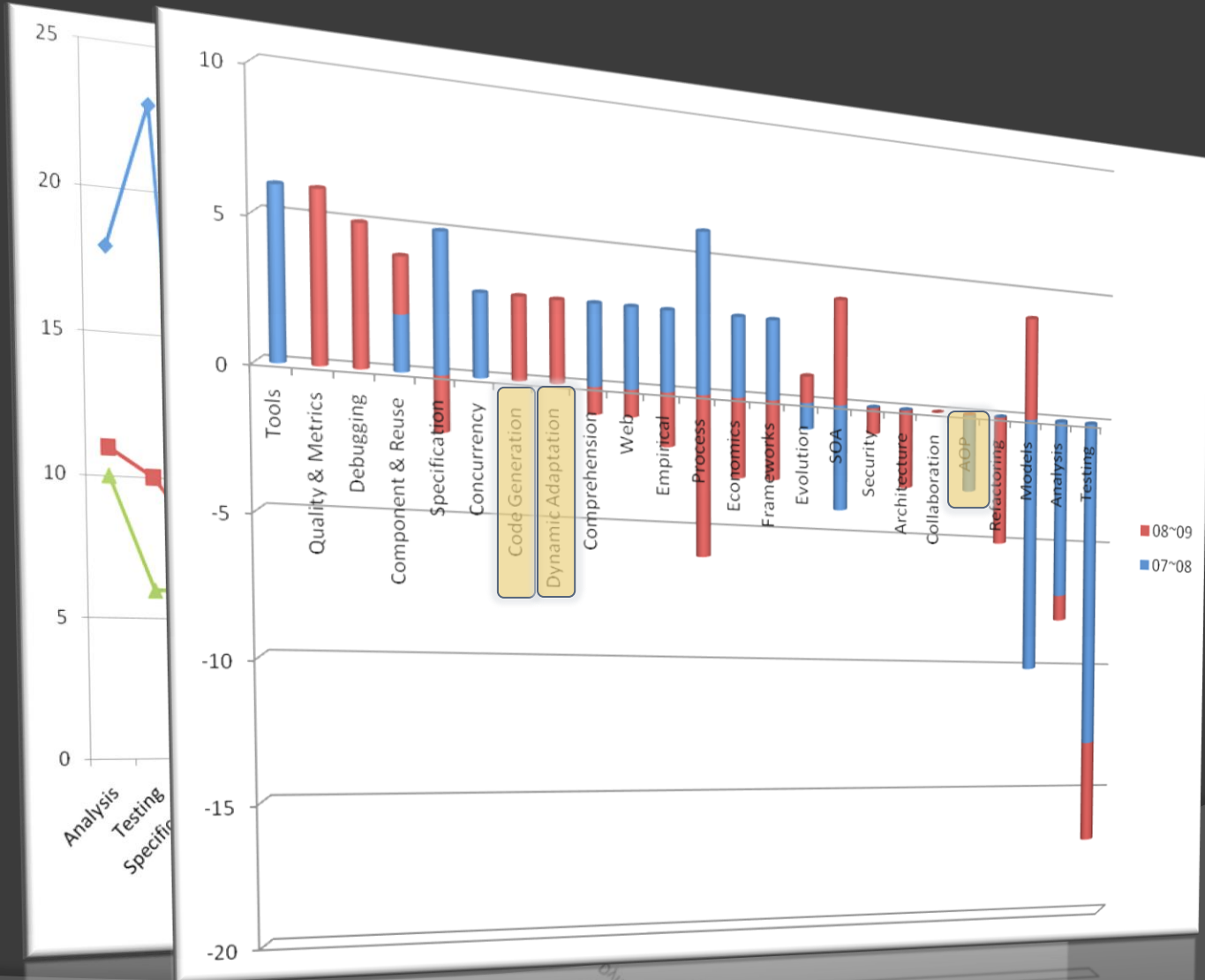
문제가 되는 점은
없을까?

있다면 어떻게
해결해야 할까?

뭘든까?

왜든왜야 뭘까?

A New Hope



Test Coverage

테스트 묶음에 의해 수행된
코드의 크기

쉽게 예상할
수 있는 크기



어렵게 예상해
야 하는 크기

In APSEC 2009

16th APSEC (Asia-pacific Software engineering Conference)

Test Coverage Metric for Two-staged Language with Abstract Interpretation

Taeksu Kim, Chunwoo Lee, Kiljoo Lee, Soohyun Baik, Kwangkeun Yi, and Chisu Wu.

December 2009.



100% decision coverage. However, the correct decision coverage value should be 67% because the example code in figure 1 will run actually a code in figure 2.

```
function foo(a, b) {  
  if (b > 0 && a > 0) return 3;  
  else if (b > 0) return 4;  
  else return 0;  
}
```

Figure 2. Actual running code of figure 1

It is needed to analyze what code fragments would be generated in execution time to guarantee the correctness of a test coverage. In this paper, we designed static analysis...

I. INTRODUCTION

I. INTRODUCTION

다음 단계는...

- ◎ Symbolic Execution
- ◎ Concolic Testing

Concolic Testing

Concolic Testing을 Multi-staged 언어
에 적용할 때

분기가 동적으로
생기긴 하지만...

실제 수행 과정은
어떻게 될까?

유기동적분기...

의동적분기...

목표

정리 1. 인자가 a 인 함수에 대해 인자에 임의의 값 a_0 을 대입해 수행한 결과 얻어낸 조건 부등식 집합을 C 라 하자. 이 때 C 를 만족하는 값의 집합 V 내의 임의의 값을 인자에 대입해 프로그램을 수행하면 항상 수행하는 경로는 같다.

정리 2.

$$\sigma_0, \emptyset, \Gamma \vdash e : v, w, \Gamma'$$

일 때, Γ' 은 실제 프로그램 e 를 수행했을 때 얻어지는 조건식 집합을 Γ_0 이라고 하면 $\Gamma_0 = \Gamma'$.

$$\frac{\sigma, \varphi, \Gamma \vdash e : [e'], [w], \Gamma' \quad \sigma, \varphi, \Gamma' \vdash e' : v, w', \Gamma''}{\sigma, \varphi, \Gamma \vdash \text{run } e : v, w \xrightarrow{\text{run}} w', \Gamma''}$$

$$\frac{\sigma, \varphi, \Gamma \vdash e : v, w, \Gamma'}{\sigma, \varphi, \Gamma \vdash \text{lift } e : [v], [w], \Gamma'}$$

감사합니다.

Applying Concolic Testing to Software Written in Multi-staged Language

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Research On Software Analysis for Error-free

ROSAEC 3rd Workshop
Computing 2010