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ERC Workshop



Analyzing Information Spreading in Complex Networks

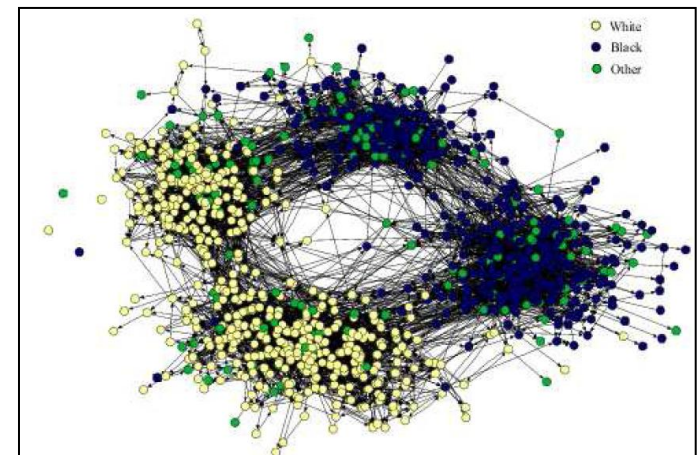
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Applied Algorithm Lab

KAIST

Information Spreading in Complex Networks

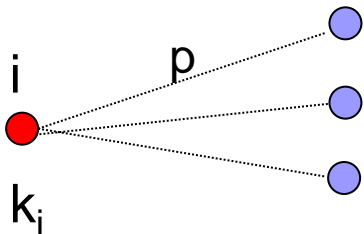
- Network models play a fundamental role as a medium for the spread of information, ideas, and influence among its members.
- Direct Marketing takes the “word-of-mouth” effects to significantly increase profits.
- Examples:
 - A company selects a small number of customers and ask them to try a new product. The company wants to choose a small group with largest influence.



Erdos-Renyi Random Graph

n nodes, connect each pair of nodes with probability p

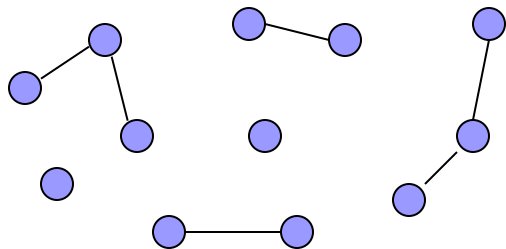
Degree dist. given by: $P_k(k_i = k) = B(n-1, p) \cong \frac{e^{-z} z^k}{k!}$ $z = pn = \langle k \rangle$



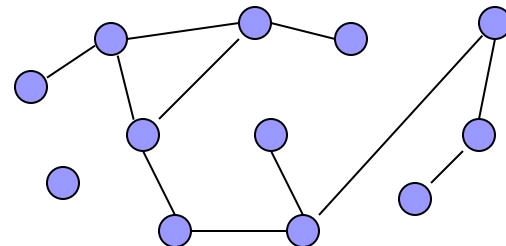
Graph structure depends only on n and p

Phase transition (For $p=z/n$, around $z=1$):

many small components



→ Single giant component



Poisson Distribution

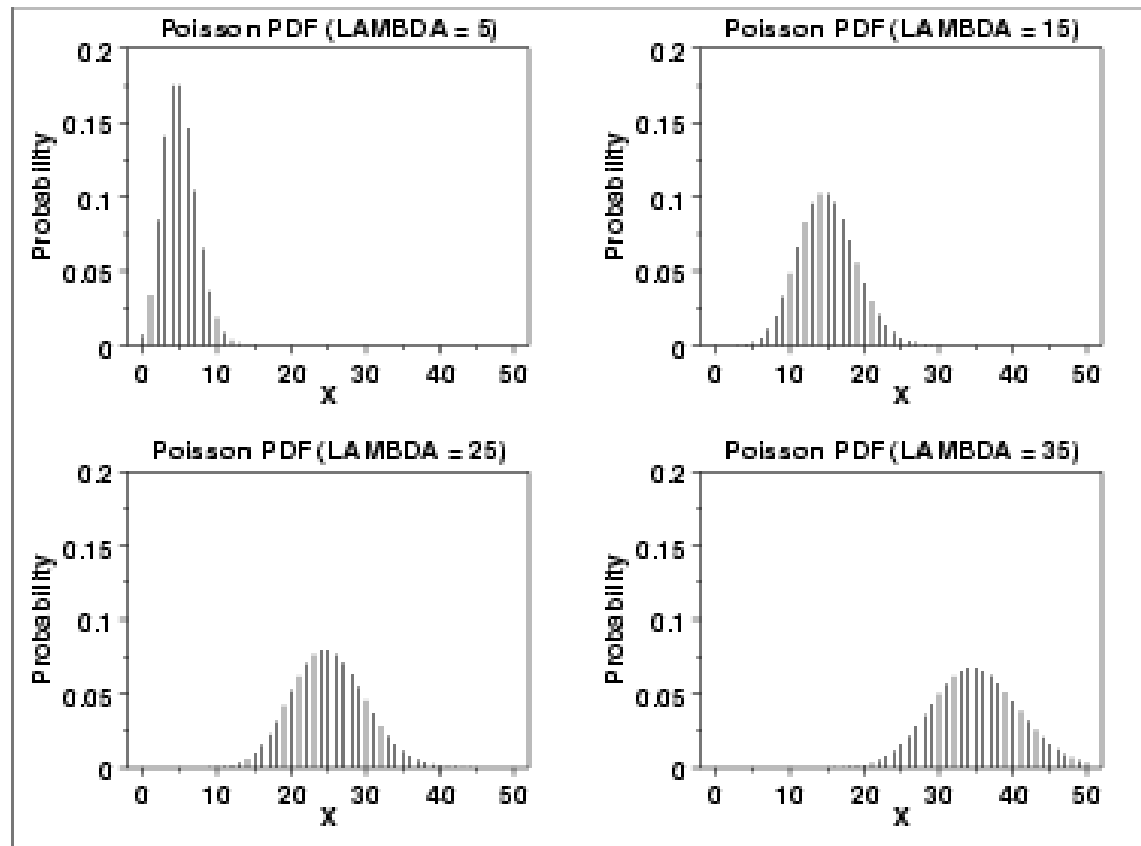
- Coming from Binomial distribution
 - Fix the expectation $z=np$
 - Let the number of trials $n \rightarrow \infty$

A Binomial distribution $B(n, z/n)$ converges to the Poisson distribution of rate z

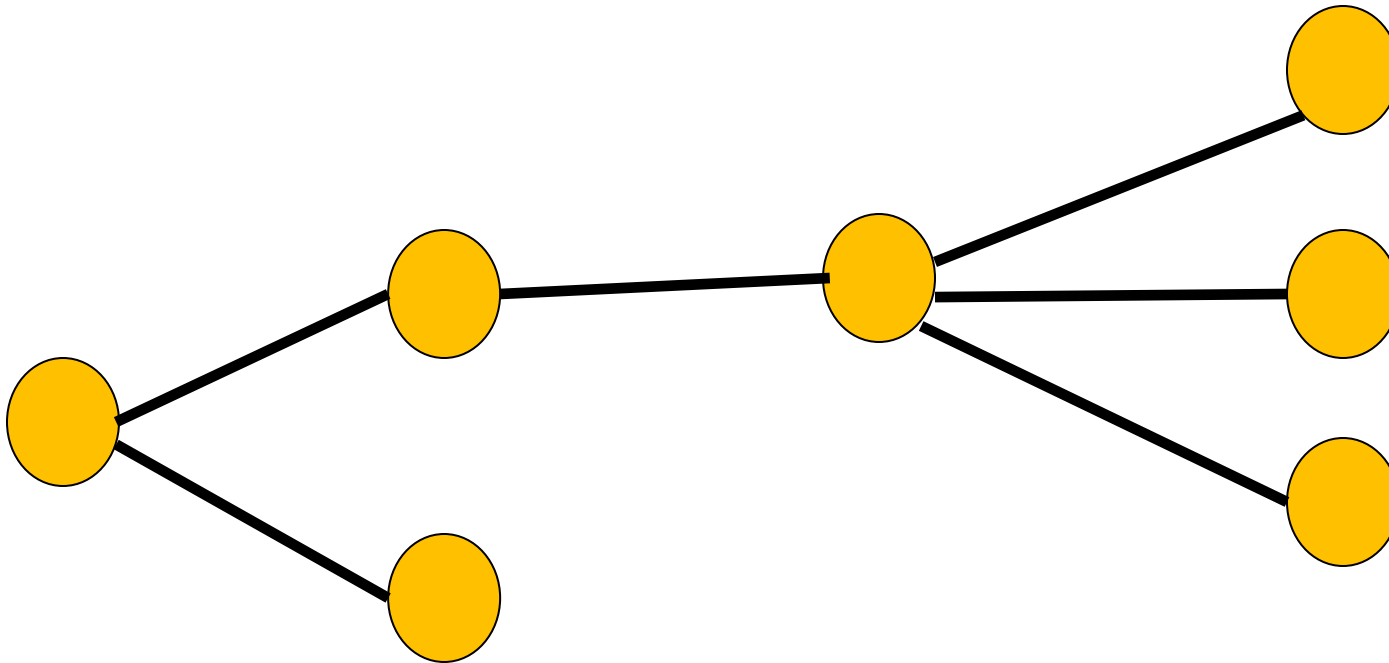
$$\Pr(X = x) = p_{\theta}(x) = \begin{cases} \frac{z^x}{x!} e^{-z} & x \geq 0 \\ 0 & \text{otherwise} \end{cases}$$

- $E[X] = z, \text{Var}(X) = z$

Plots of Poisson Distribution



Branching Process



Z is the average degree

Giant component size

u ~ prob. node is not in giant component

S ~ fraction of nodes in giant component

$$S = 1 - u$$

If node is not in giant component, then neither are its **neighbors**:

$$u = \sum P_k u^k = e^{-z} \sum \frac{(uz)^k}{k!} = e^{z(u-1)}$$

Notice: this depends crucially on degree dist.

Fraction of nodes in giant component:

$$S = 1 - e^{-zS}$$

We can also calculate the solution iteratively.

Giant component size

$$S = 1 - e^{-zS}$$

- For $z < 1$, the only non-negative solution is $S = 0$.
- For $z > 1$ (after the phase transition), the only non-negative solution is the **fractional size of the giant component**.

Scale Free Networks

One particularly **ubiquitous** degree distribution form is the **power law**:

$$P(k) \sim k^{-\gamma}$$

Network	Size	γ^{in}	γ^{out}	Ref
www	2×10^8	2.1	2.71	Broder (2000)
Movie actors	2.12×10^5	2.3	2.3	(Barabasi 1999)
Word co-occurrence	4.62×10^5	2.7	2.7	(Cancho 2001)

So what underlying mechanism is responsible for the power law distribution?

There is something **special** about power law distributions...

Power Laws and Scale-Free behavior

Scale-free Criteria: $p(ax) = f(a)p(x)$

Differentiating above w.r.t. a and considering the cases $x=1, a=1$ yields:

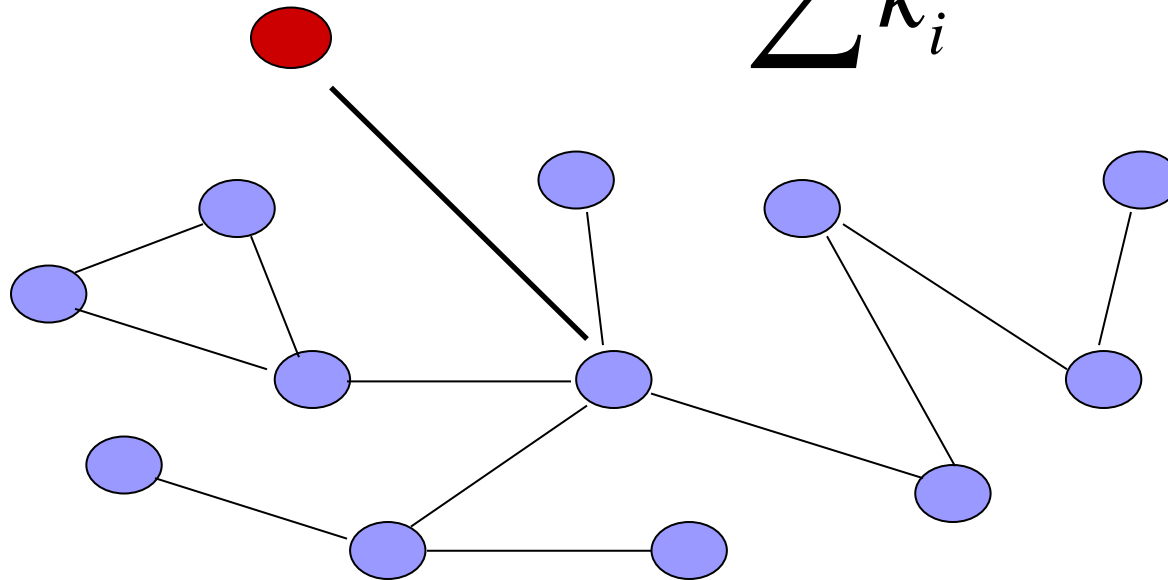
$$x \frac{dp}{dx} = \frac{p'(1)}{p(1)} p(x) \quad \text{with the solution:} \quad p(x) \propto x^{-\alpha}$$

Power law distributions are the **only** functions that satisfy the **Scale Free Criteria**.

Preferential Attachment Model

Attach new node to existing graph with probability:

$$\Pi_i = \frac{k_i}{\sum k_i}$$



This model was derived in the 1950's by Herbert Simon.

- who won a Nobel Prize in economics for entirely different work.

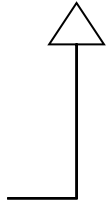
Preferential Attachment Model

$N(k, t) \sim$ number of nodes with degree k at time t

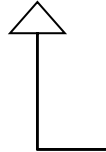
A decrease in $N(k, t)$ implies an increase in $N(k + 1, t + 1)$

$$p(k, t + 1) = \left(\frac{k - 1}{2t} \right) p(k - 1, t) + \left(1 - \frac{k}{2t} \right) p(k, t)$$

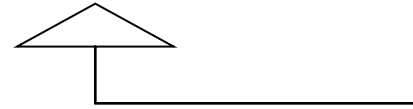
Prob. that a node born at time t_i has degree k at time $t+1$



Prob. a new node pref. attaches to a node with degree $k-1$



Prob. a new node does not pref. attach to a node with degree k



The above becomes:

$$P(k) = \begin{cases} (k - 1) / (k + 2) * P(k - 1) & k \geq m + 1 \\ 2 / (m + 2) & k = m \end{cases}$$

The degree dist. is:

$$P(k) = \lim_{t \rightarrow \infty} \left(\sum_{t_i} p(k, t) \right) / t$$

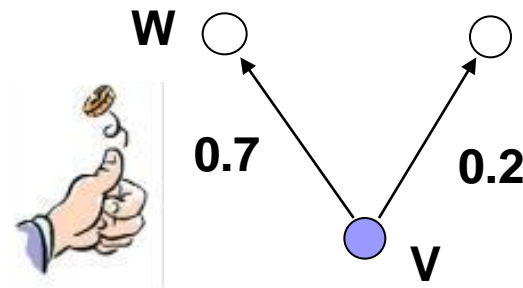
$$P(k) = \frac{2m(m + 1)}{k(k + 1)(k + 2)}$$

Information Spreading Models

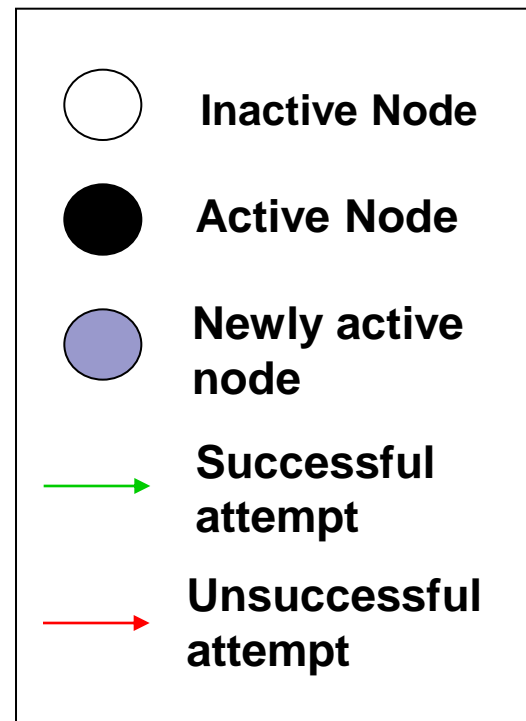
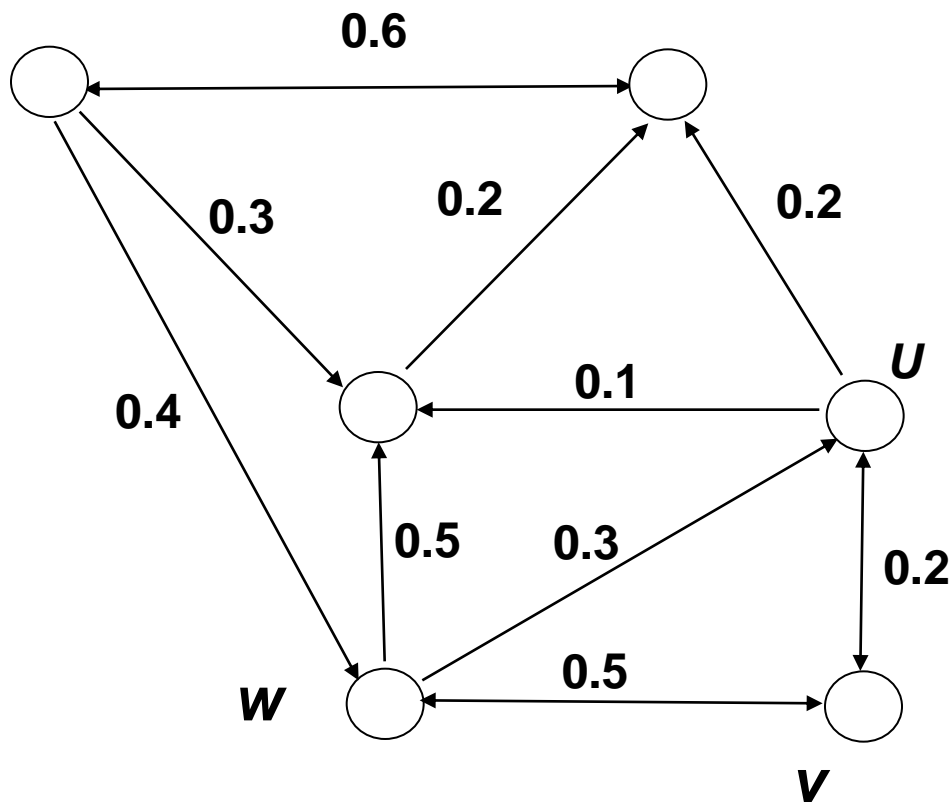
- A network is represented as a graph. Each customer is considered as a node.
- Each node can be either active (accept the information) or inactive.
- By the “word-of-mouth” effects, each node’s tendency to become active increases monotonically as more of its neighbors become active.
- Assumption: node can switch to **active** from **inactive**, but does not switch in the other direction.

Independent Cascade Model

- Starts with an initial set of active nodes A_0
- The diffusion process keeps in discrete steps
 - When node V first becomes active in step t , it is given a single chance to activate each currently inactive neighbor W . It succeeds at probability $p_{V,W}$ – a parameter of the system.



Example

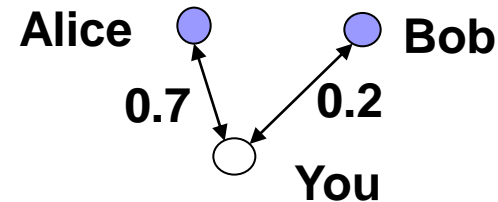


Stop!

Linear Threshold Model

- A node v is influenced by each neighbor w according to a *weight* $b_{v,w}$ such that

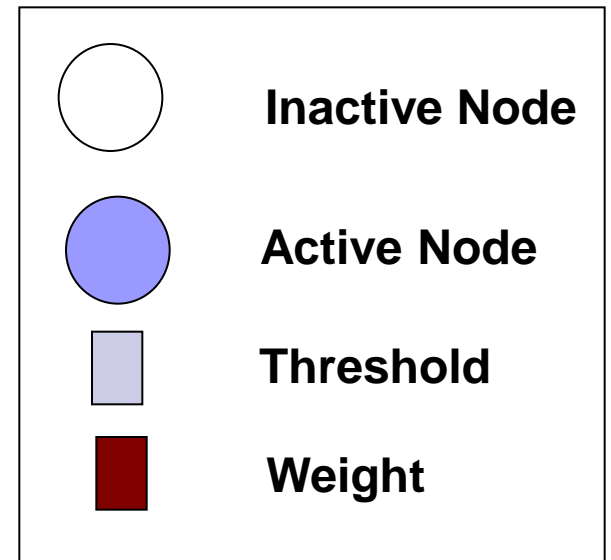
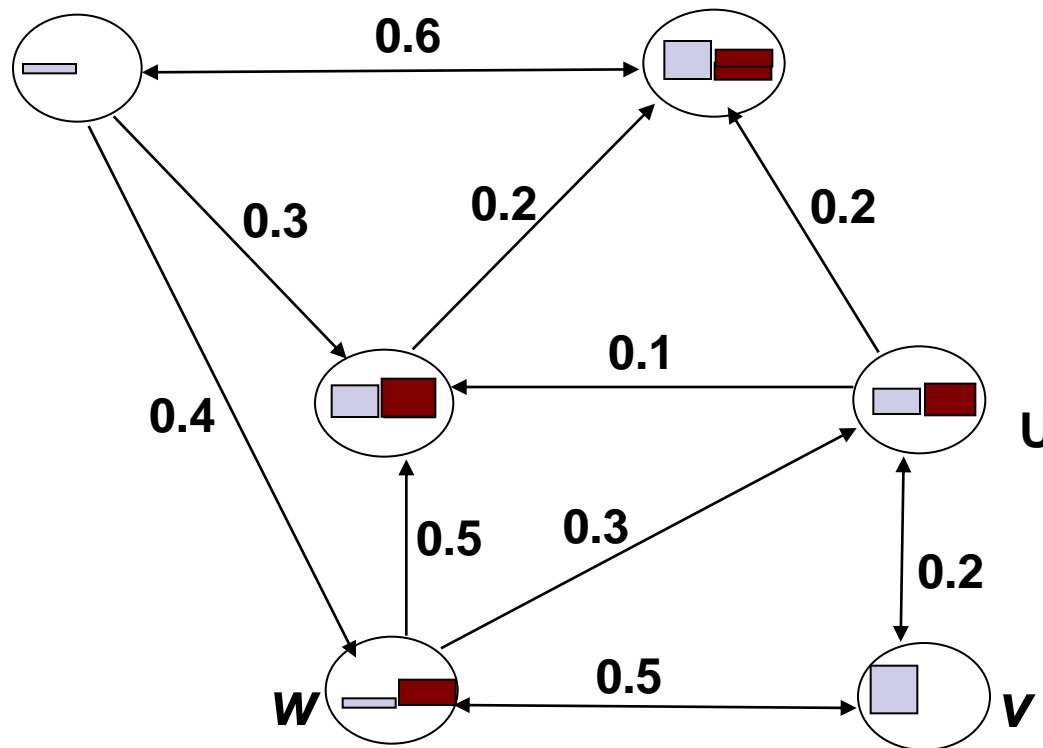
$$\sum_{w \text{ neighbor of } v} b_{v,w} \leq 1$$



- Each node v has a threshold θ_v which is chosen from the interval $[0,1]$.
- A node v becomes active if

$$\sum_{\substack{w \text{ neighbor of } v \\ w \text{ is active}}} b_{v,w} \geq \theta_v$$

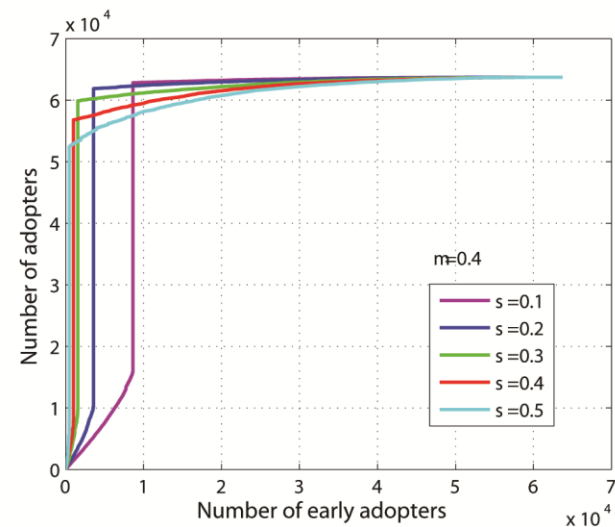
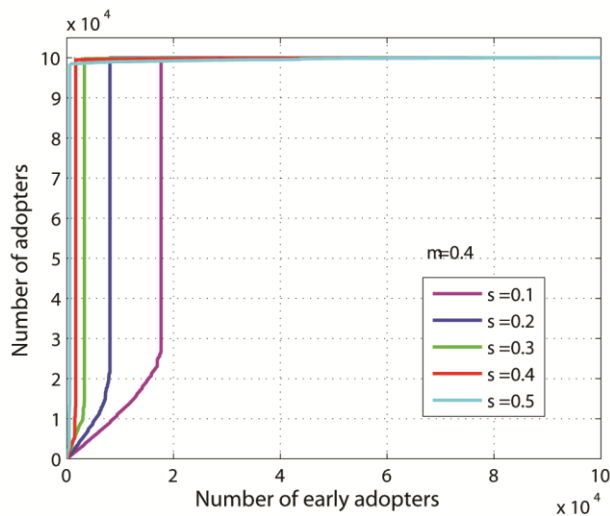
Example



Stop!

Threshold phenomenon in information spreading

- We observe **giant information spreading** suddenly occurs according to a parameter in real social networks.
- We are developing models, and mathematically and statistically analyzing this phenomenon.



Influence Maximization Problem

- Define $\sigma(A)$ to be the number of active nodes at the end of the information spreading process.
- Problem Definition:
 - Given a parameter k , find a k -node set A to maximize $\sigma(A)$.
 - NP-hard for both independent cascade model and linear threshold model.
- We propose a novel recursive algorithm that approximately computes the influence.
 - Our algorithm empirically performs very well.

Applied Algorithm Lab

■ Graduate Students

- Nam-ju Kwak (branching process in social networks)
- Boyoung Kim (decentralized ranking learning)
- Yongsub Lim (graphical model, multi-agent system analysis)
- Sungsu Lim (modeling random scale-free network)
- Wooram Heo (influence maximization)
- Seulki Lee (threshold analysis of information spreading)

■ Undergraduate Research Program

- Taejin Chin (sorting algorithm for partially sorted list)