

SAT Algorithms

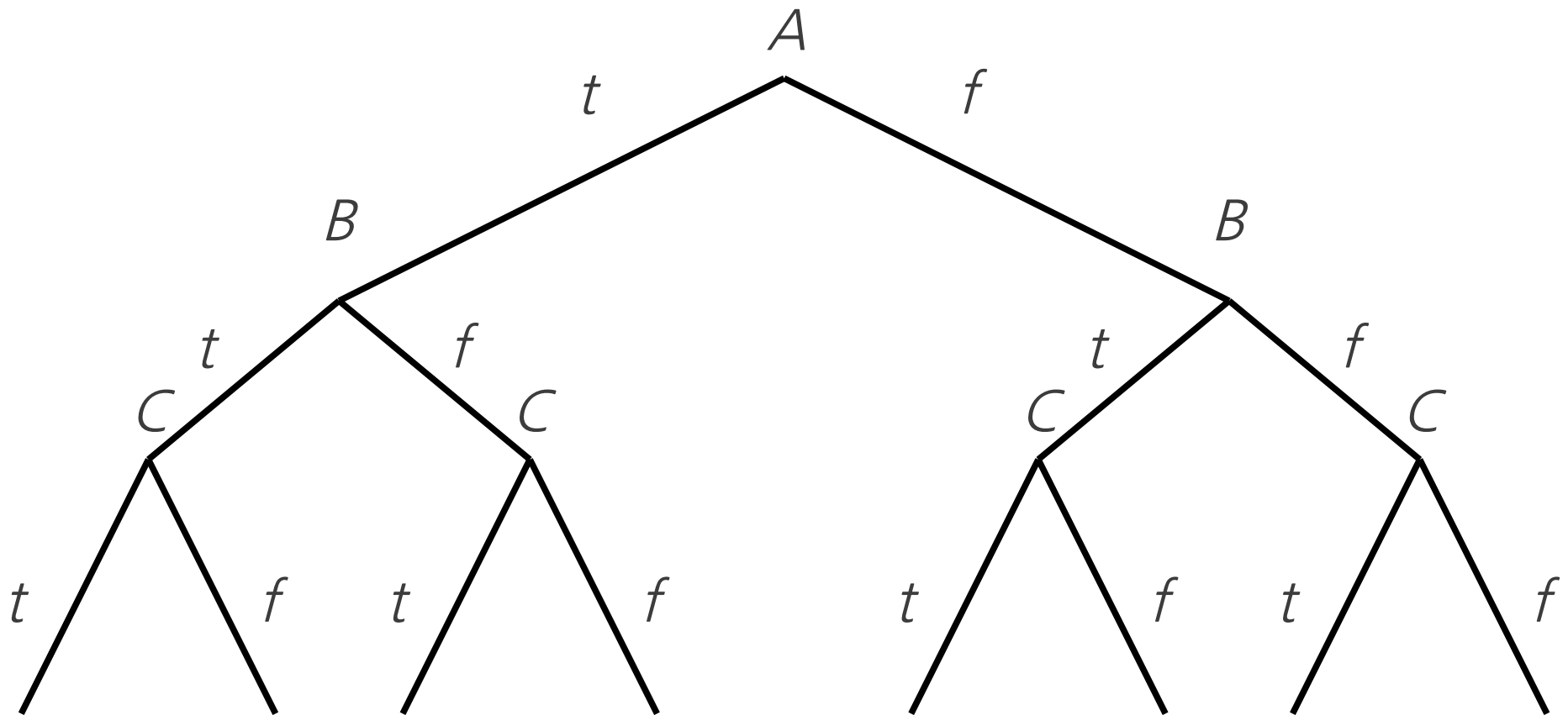
DPLL and CDCL

서울대학교 프로그래밍 연구실
조성근

$$(\neg A \vee B) \wedge (\neg B \vee C)$$

이 식을 참으로 만드는 A, B, C가 있나?

$$(\neg A \vee B) \wedge (\neg B \vee C)$$



DPLL 알고리즘

(정의 (dpll 제약조건)

(조건

[(다풀었니? 제약조건) (공집합)]

[(충돌이있니? 제약조건) (못풀어)]

[(다르니? (dpll (x를참으로 제약조건)) (못풀어))
(추가 (결과) (x=참))]

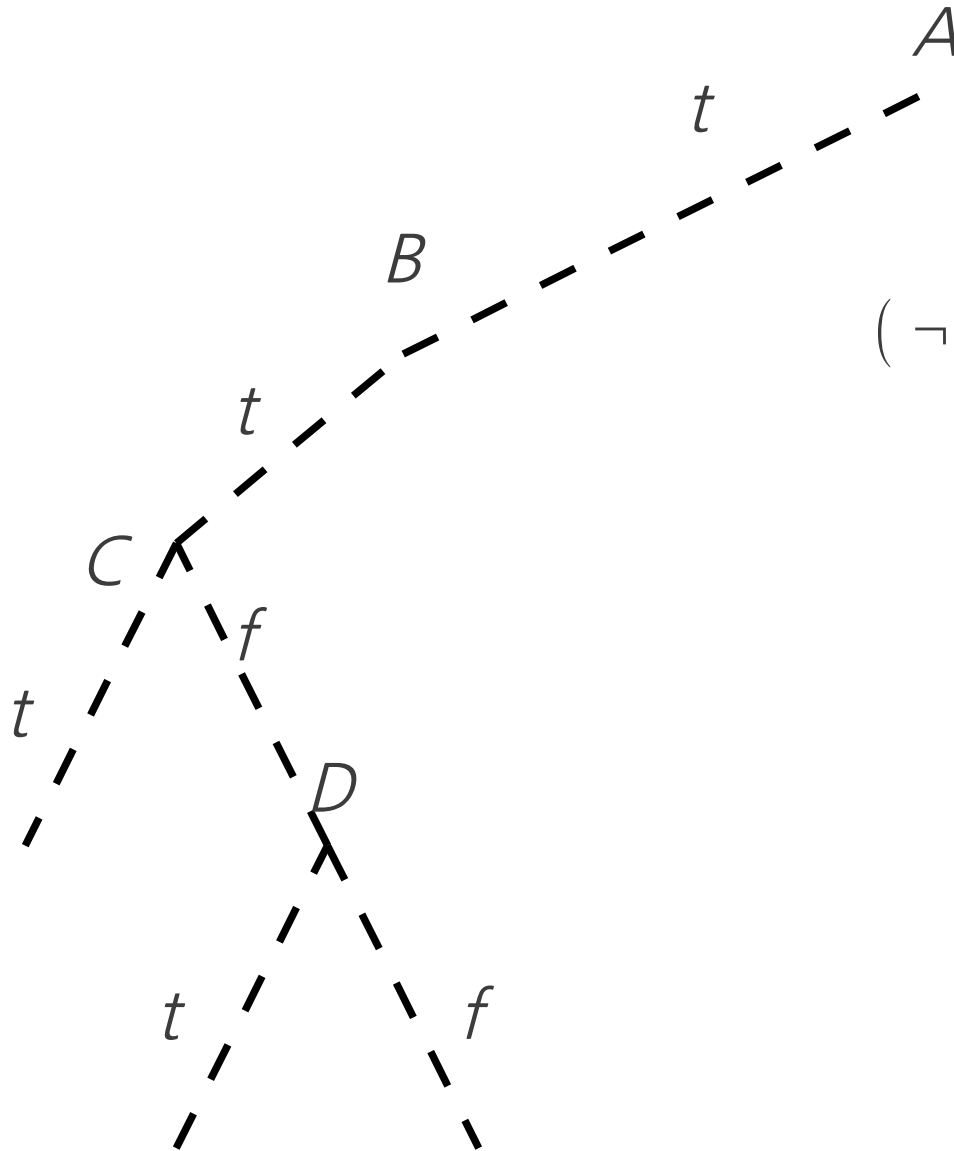
[(다르니? (dpll (x를거짓으로 제약조건)) (못풀어))
(추가 (결과) (x=거짓))]

[(참) (못풀어)]

)

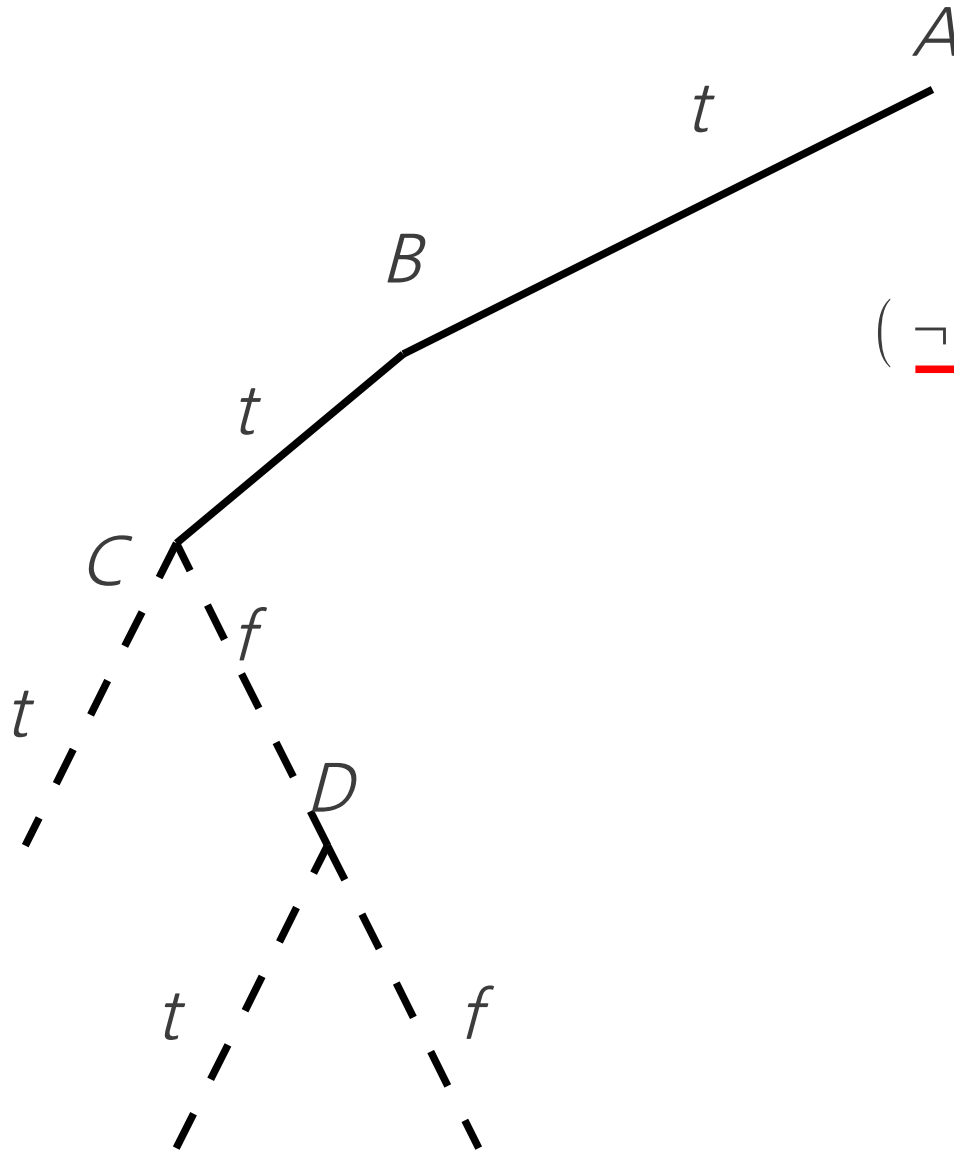
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DPLL



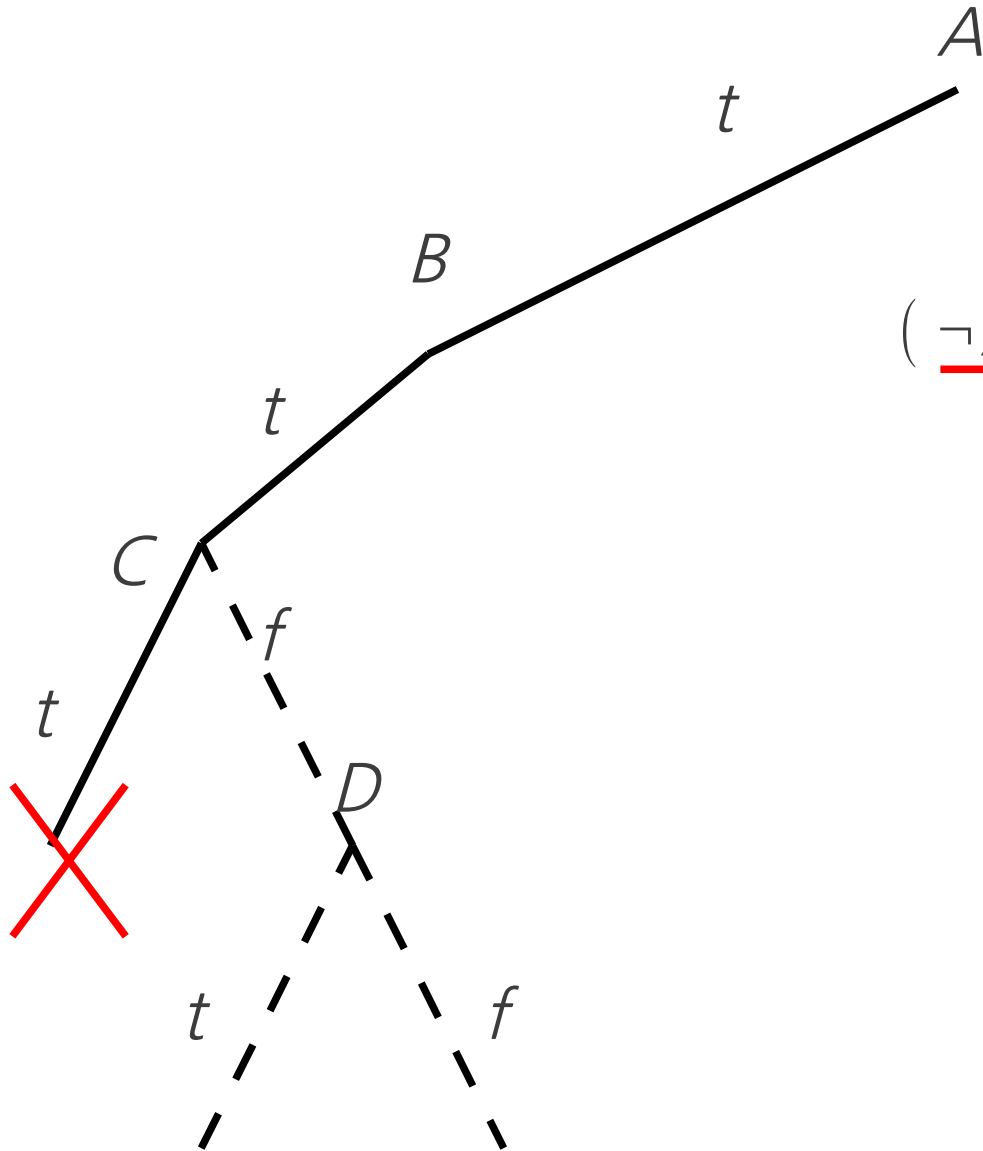
$$(\neg A \vee B) \wedge (\neg B \vee \neg C) \wedge (C \vee \neg D)$$

DPLL



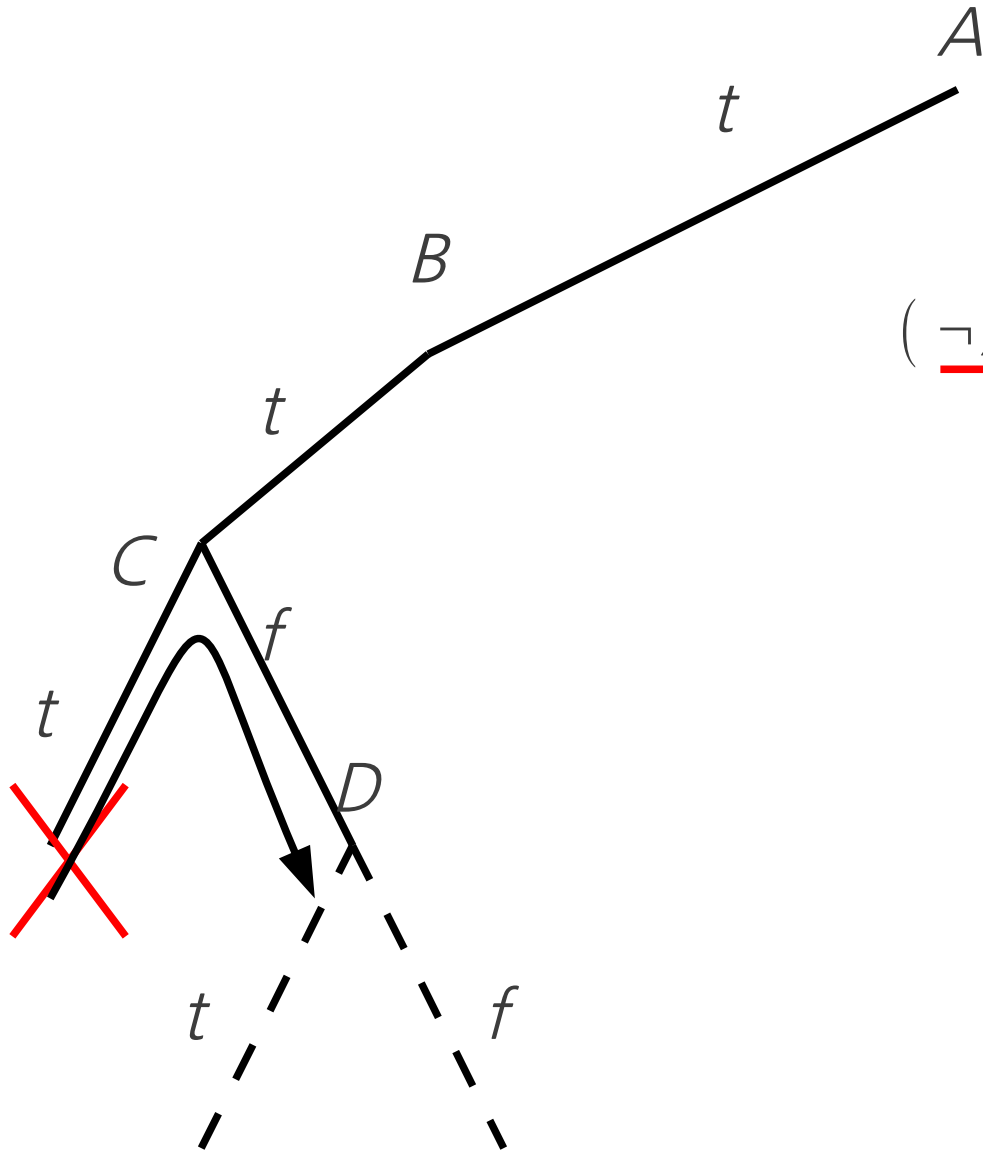
$$(\underline{\neg A} \vee \underline{B}) \wedge (\underline{\neg B} \vee \neg C) \wedge (C \vee \neg D)$$

DPLL



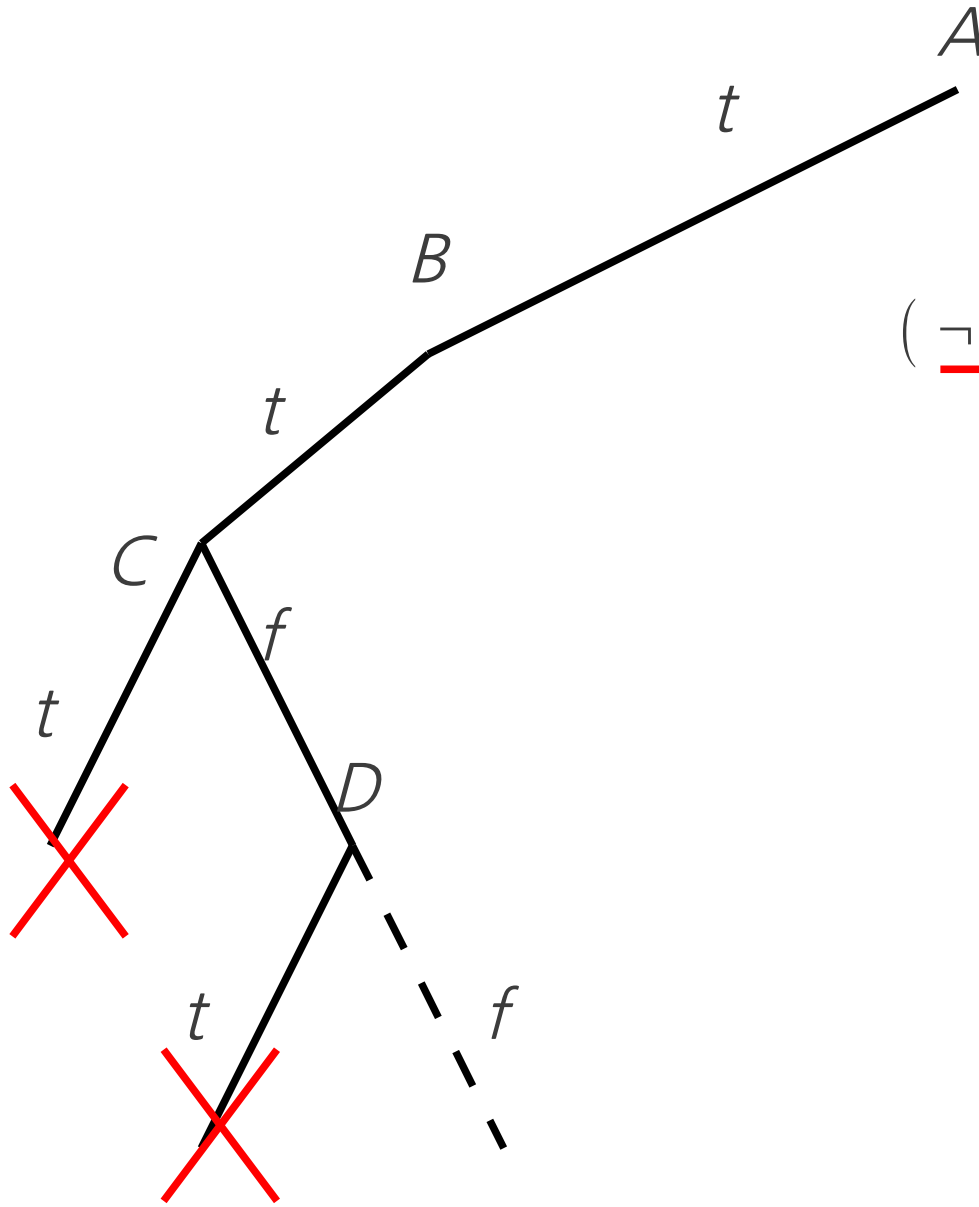
$$(\underline{\neg A} \vee \underline{B}) \wedge (\underline{\neg B} \vee \underline{\neg C}) \wedge (\underline{C} \vee \neg D)$$

DPLL



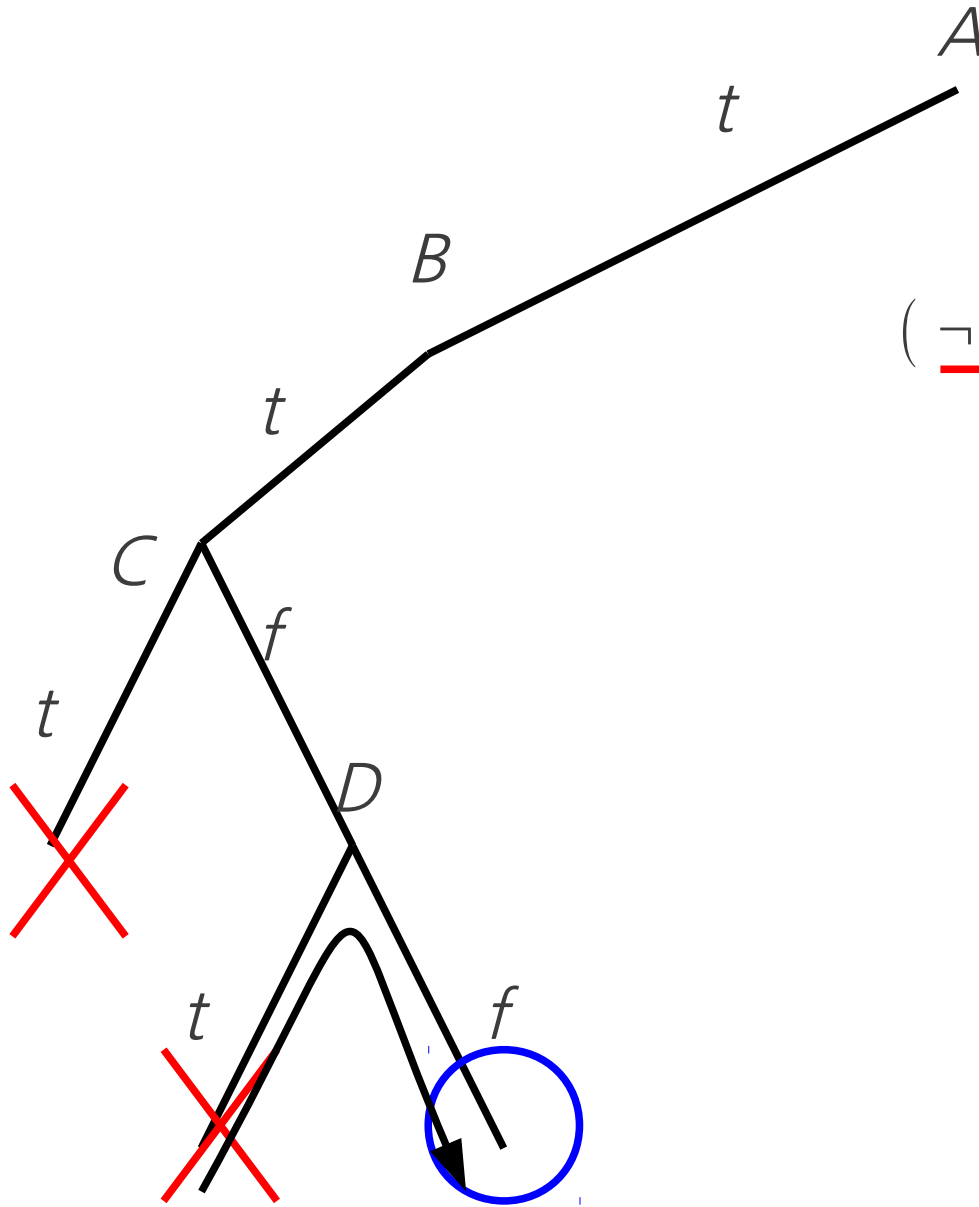
$$(\underline{\neg A} \vee \underline{B}) \wedge (\underline{\neg B} \vee \underline{\neg C}) \wedge (\underline{C} \vee \neg D)$$

DPLL



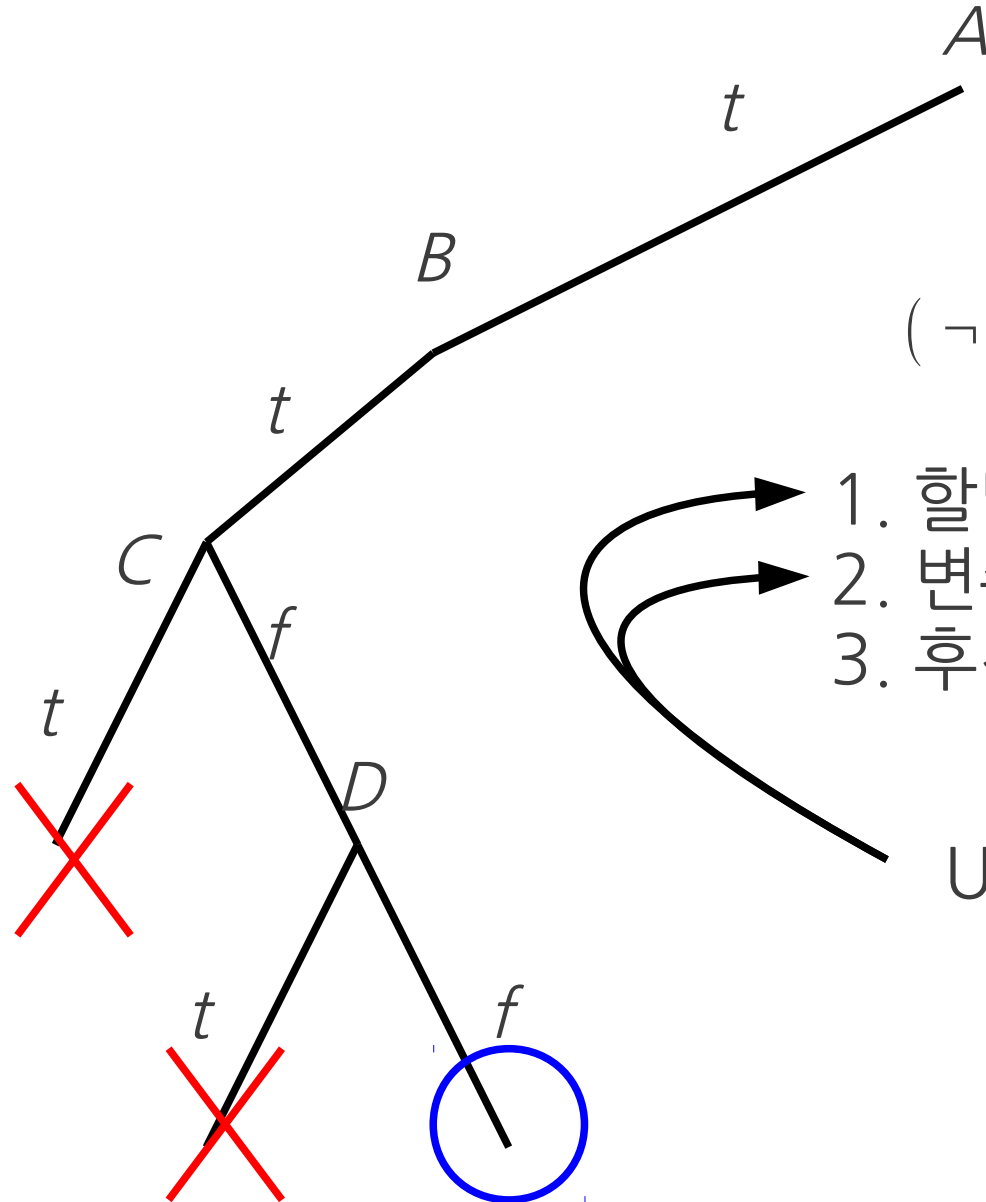
$$(\underline{\neg A} \vee \underline{B}) \wedge (\underline{\neg B} \vee \underline{\neg C}) \wedge (\underline{C} \vee \underline{\neg D})$$

DPLL



$$(\underline{\neg A} \vee \underline{B}) \wedge (\underline{\neg B} \vee \underline{\neg C}) \wedge (\underline{C} \vee \underline{\neg D})$$

DPLL



$$(\neg A \vee B) \wedge (\neg B \vee \neg C) \wedge (C \vee \neg D)$$

1. 할당하는 변수의 순서
2. 변수에 좋은 값 주기
3. 후진(backtracking) 잘 하기

Unit Resolution

Unit Resolution

- literal이 하나 밖에 없는 clause가 있을 때
- 예
 - $(\underline{B}) \wedge (\neg B \vee \neg C) \wedge (C \wedge \neg D)$
 - $(\neg A \vee B) \wedge (\underline{\neg C}) \wedge (C \wedge \neg D)$
- 그 literal이 참이 되도록 값을 할당한다.

CDCL

- Conflict-Driven Clause Learning
- DPLL로부터 몇가지 기능이 추가된 버전
 - 충돌로부터 새로운 clause를 배움
 - 더 효율적인 데이터 구조(lazy data structure)
 - 주기적으로 처음부터 다시 시작
 - 새롭게 배운 clause가 너무 많아지면 조금 지우기

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Implication Graph

어떤 변수에게 참/거짓을 할당하였더니
다른 변수가 참/거짓이어야 하더라.

Implication Graph

$$\begin{aligned}\varphi &= \omega_1 \wedge \omega_2 \wedge \omega_3 \wedge \omega_4 \wedge \omega_5 \wedge \omega_6 \\ &= (x_1 \vee x_8 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (x_2 \vee x_3 \vee x_4) \wedge \\ &\quad (\neg x_4 \vee \neg x_5) \wedge (x_7 \vee \neg x_4 \vee \neg x_6) \wedge (x_5 \vee x_6)\end{aligned}$$

값을 할당하는 변수의 순서는

$$x_7, x_8, x_1, \dots$$

Implication Graph

$$\begin{aligned}\varphi &= \omega_1 \wedge \omega_2 \wedge \omega_3 \wedge \omega_4 \wedge \omega_5 \wedge \omega_6 \\ &= (x_1 \vee x_8 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (x_2 \vee x_3 \vee x_4) \wedge \\ &\quad (\neg x_4 \vee \neg x_5) \wedge \underline{(x_7 \vee \neg x_4 \vee \neg x_6)} \wedge (x_5 \vee x_6)\end{aligned}$$

$$x_7 = 0$$

Implication Graph

$$\begin{aligned}\varphi &= \omega_1 \wedge \omega_2 \wedge \omega_3 \wedge \omega_4 \wedge \omega_5 \wedge \omega_6 \\ &= (x_1 \vee \underline{x_8} \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (x_2 \vee x_3 \vee x_4) \wedge \\ &\quad (\neg x_4 \vee \neg x_5) \wedge (\underline{x_7} \vee \neg x_4 \vee \neg x_6) \wedge (x_5 \vee x_6)\end{aligned}$$

$$x_8 = 0$$

$$x_7 = 0$$

Implication Graph

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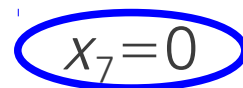
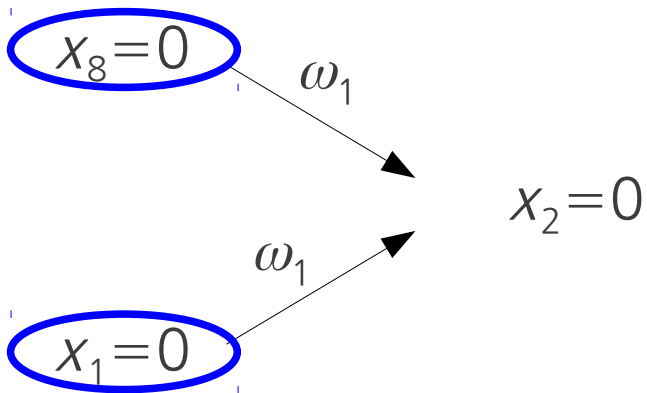
$$x_8 = 0$$

$$x_1 = 0$$

$$x_7 = 0$$

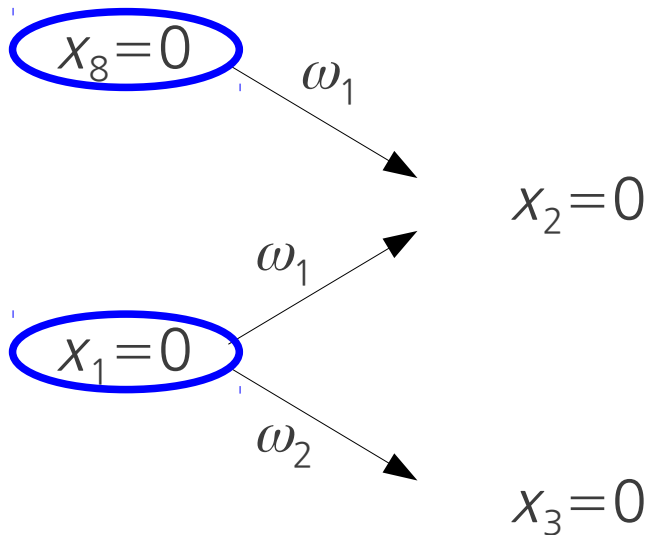
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Implication Graph

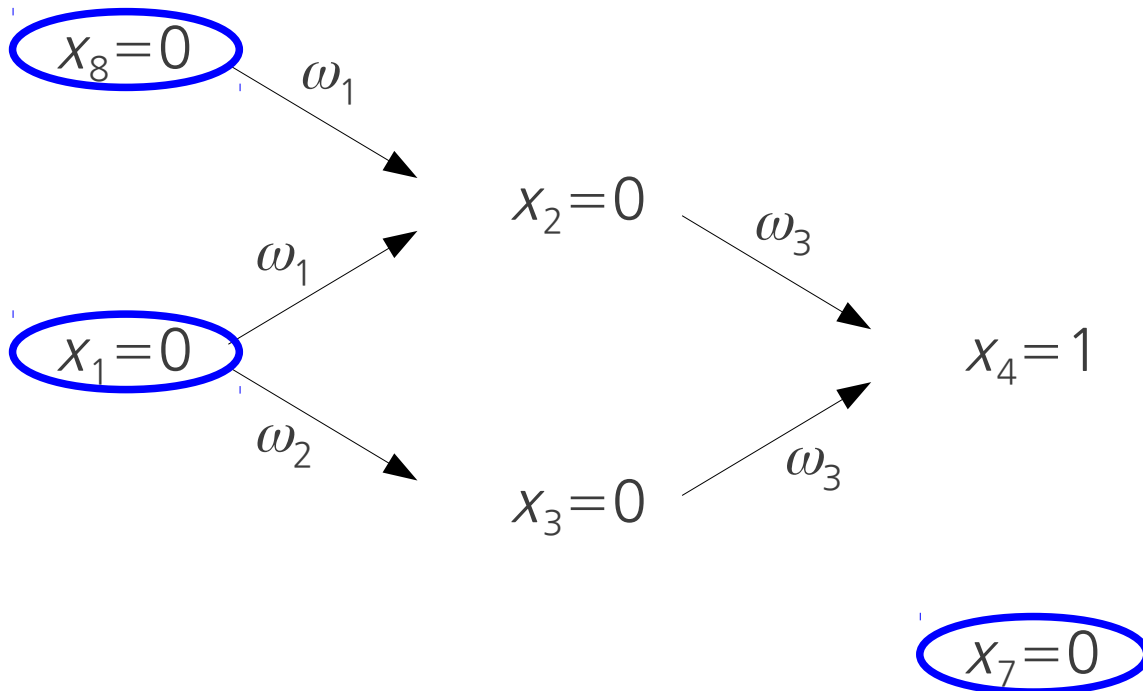
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$$x_7=0$$

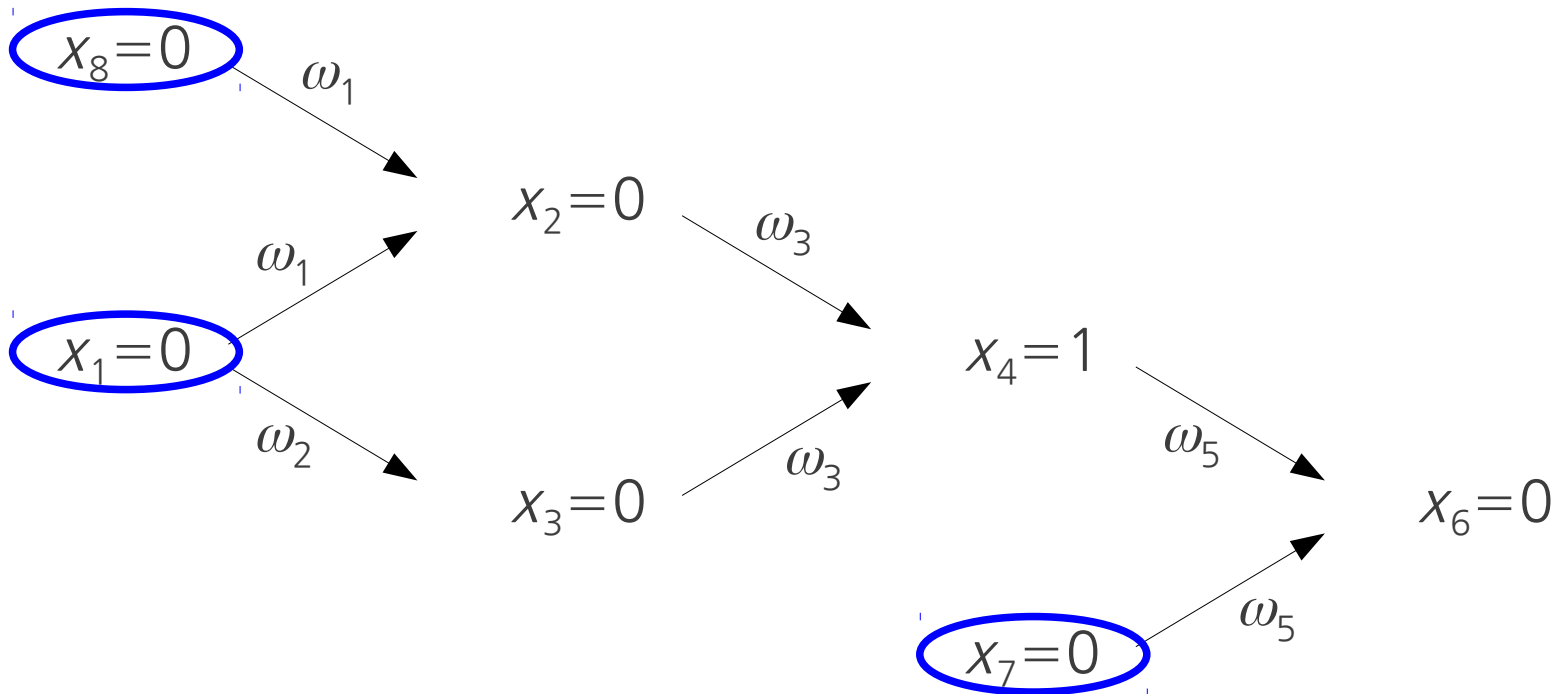
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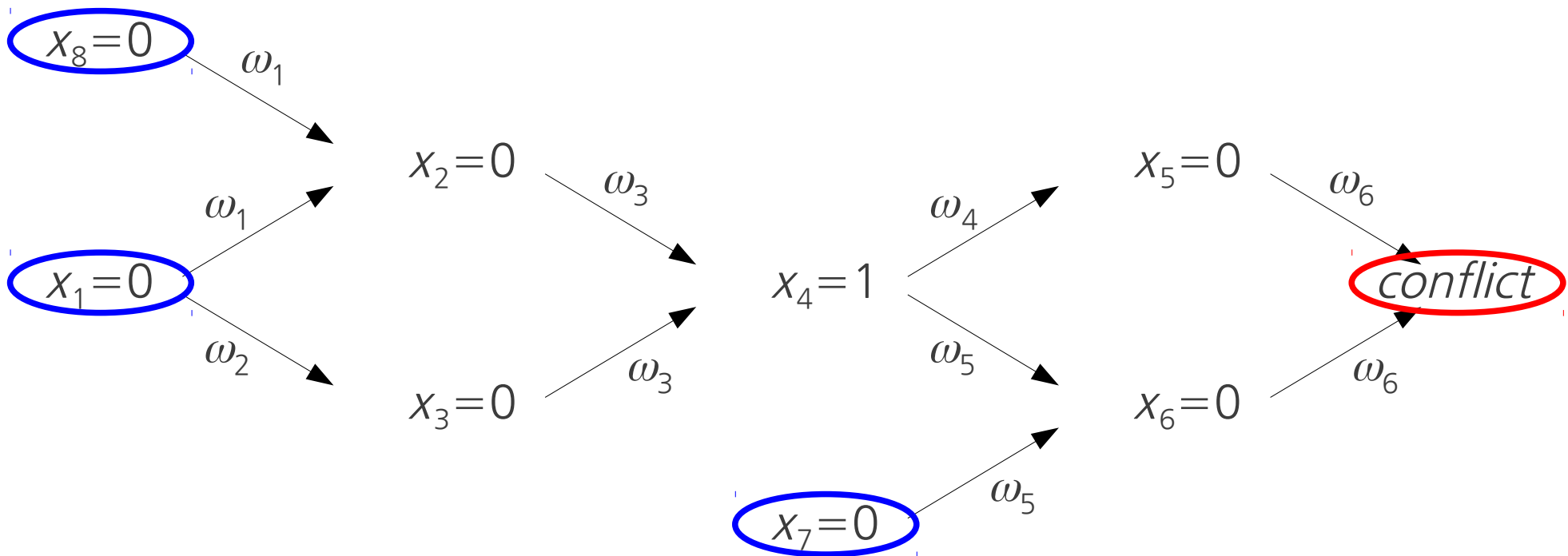
Implication Graph

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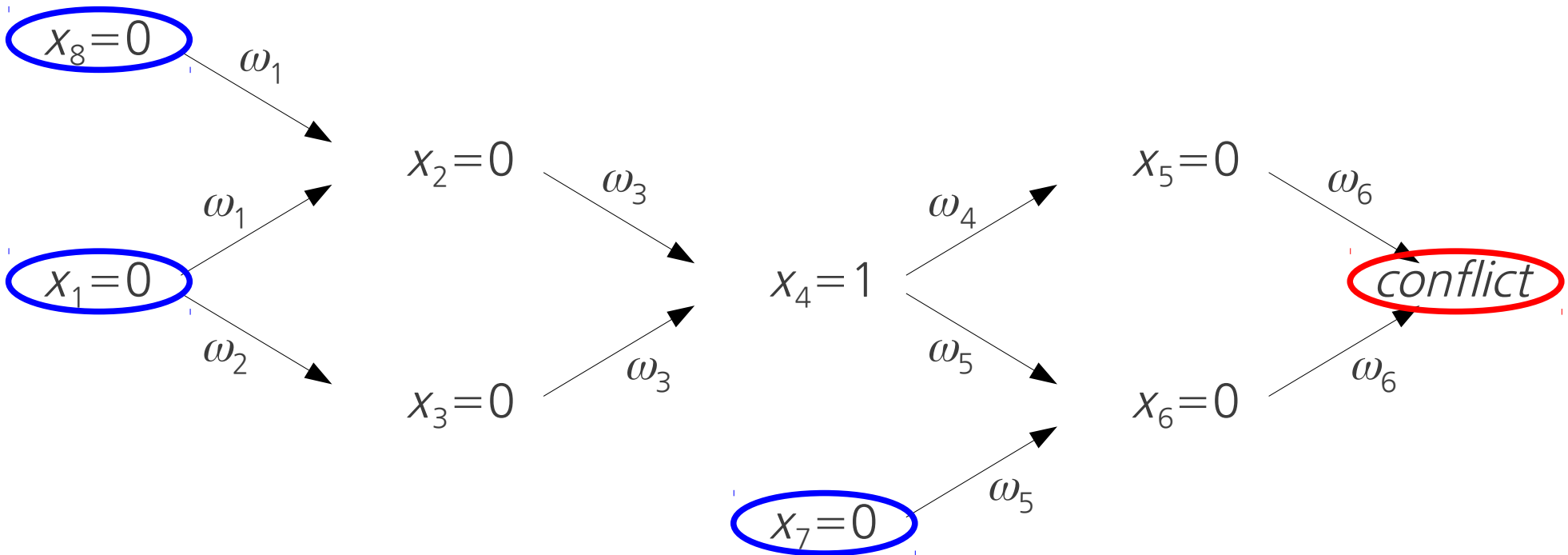
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Clause 배우기

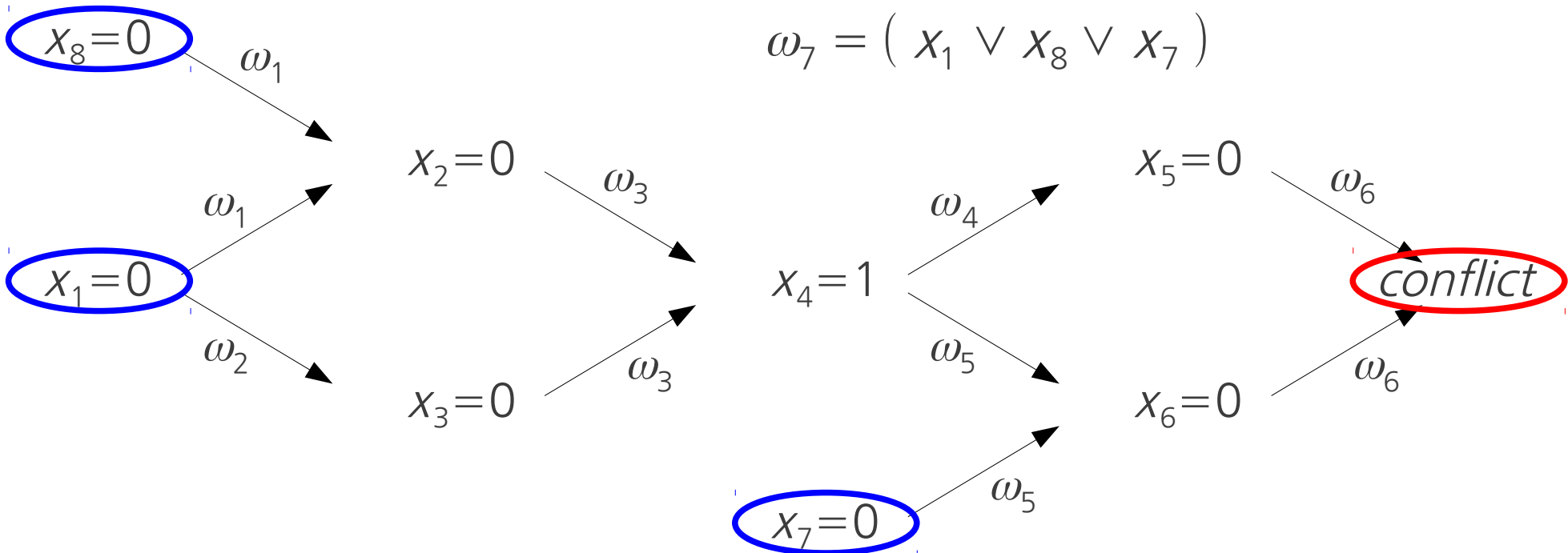
- 충돌을 일으키는데 기여하는 clause들로부터 새로운 clause를 만들어 배우자.



Clause 배우기

$$\begin{aligned}
 \omega_6 &= (x_5 \vee x_6) && (\neg x_4 \vee x_6) \\
 \omega_4 &= (\neg x_4 \vee \neg x_5) && (\neg x_4 \vee x_7) \\
 \omega_5 &= (x_7 \vee \neg x_4 \vee \neg x_6) && (x_2 \vee x_3 \vee x_7) \\
 \omega_3 &= (x_2 \vee x_3 \vee x_4) && (x_1 \vee x_8 \vee x_3 \vee x_7) \\
 \omega_1 &= (x_1 \vee x_8 \vee \neg x_2) && (x_1 \vee x_8 \vee x_3 \vee x_7) \\
 \omega_2 &= (x_1 \vee \neg x_3) &&
 \end{aligned}$$

$$\omega_7 = (x_1 \vee x_8 \vee x_7)$$



Clause 배웠다

$$\begin{aligned}\varphi &= \omega_1 \wedge \omega_2 \wedge \omega_3 \wedge \omega_4 \wedge \omega_5 \wedge \omega_6 \\ &= (x_1 \vee x_8 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (x_2 \vee x_3 \vee x_4) \wedge \\ &\quad (\neg x_4 \vee \neg x_5) \wedge (x_7 \vee \neg x_4 \vee \neg x_6) \wedge (x_5 \vee x_6)\end{aligned}$$

$$\omega_7 = (x_1 \vee x_8 \vee x_7)$$

$$\begin{aligned}\varphi &= \omega_1 \wedge \omega_2 \wedge \omega_3 \wedge \omega_4 \wedge \omega_5 \wedge \omega_6 \wedge \omega_7 \\ &= (x_1 \vee x_8 \vee \neg x_2) \wedge (x_1 \vee \neg x_3) \wedge (x_2 \vee x_3 \vee x_4) \wedge \\ &\quad (\neg x_4 \vee \neg x_5) \wedge (x_7 \vee \neg x_4 \vee \neg x_6) \wedge (x_5 \vee x_6) \wedge \\ &\quad (x_1 \vee x_8 \vee x_7)\end{aligned}$$

결론

- 간단한 DPLL
 - 후진(backtracking)에 기초한 SAT 알고리즘
- CDCL
 - DPLL의 업그레이드 버전
 - 충돌로부터 제약조건을 배우는 SAT 알고리즘

감사합니다.