An Empirical Study on the Rewritability of the with Statement in JavaScript

ERC Workshop (16 January to 19 January, 2012)

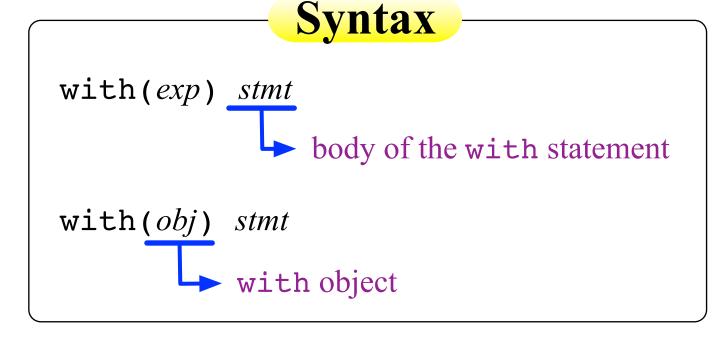
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Motivation

- Provide empirical data for with statements used in real-world JavaScript applications.
- Amount
- Usage patterns
- Check rewritability of with statements in real world.
- Rewriting rules to remove with statements by replacing them with other statements

Introduction of the with statement

• Syntax and semantics in ECMAScript



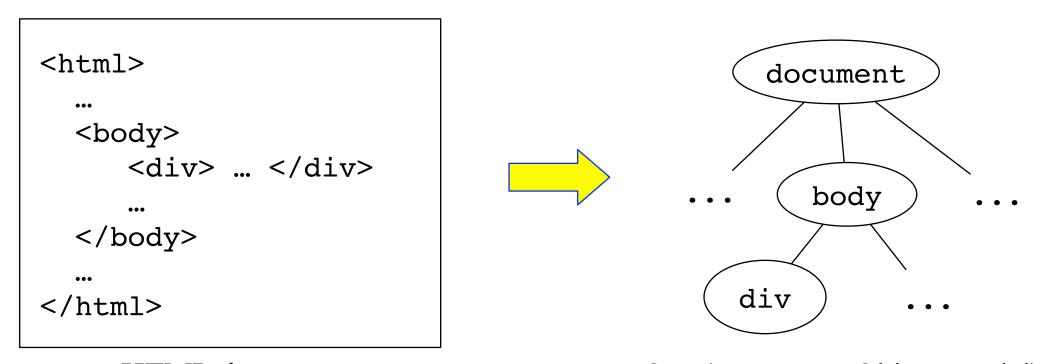
Semantics

- 1. Evaluate *exp* to a JavaScript object
- 2. Add *obj* to to the front of the scope chain
- 3. Evaluate *stmt*
- 4. Remove *obj* from the scope chain

The with statement introduces a new scope at run time.

Good Parts

• Provide a convenient way to develop dynamically changing web contents



HTML document

DOM (Document Object Model) tree

Accessing Contents in the div Tag

document.body.children[0].style.textAlign="center"; document.body.children[0].style.fontSize=50;

with(document.body.children[0].style){

textAlign="center"; fontSize=50;

Instead of using long object accesses multiple times, the with statement enables users to use just field names in the body of the with statement.

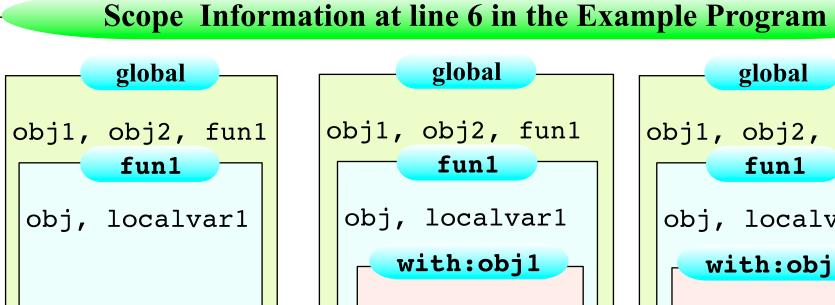
Bad Parts

- Incur performance overheads
- Make static analysis infeasible

1: var obj1={one:1}; 2: var obj2={two:2}; 3: function funl(obj){ 4: var localvar1=0; 5: with(obj) localvar1=one; 7: }

8: fun1(obj1);

9: fun1(obj2);



Before program execution At the first fun1 call

obj1, obj2, fun1 fun1 obj, localvar1 with:obj2 two

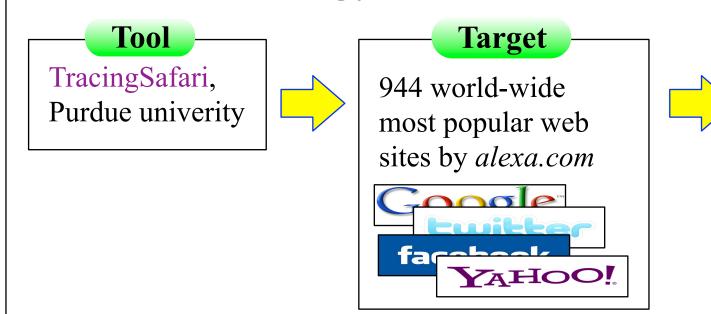
global

At the second fun1 call

Dynamic scope introduction by the with statement makes identifier resolution unclear at static time.

Real-World Usage Patterns

Methodology



• LOADING set - JavaScript code collected for 1 minute at loading time of each site

Data set

• INTERACTION set - JavaScript code collected for 2 minutes with user interactions (mouse click events) with identification

- Static with statements
- Static with keyword detected by the Rhino parser
- Executed with statements
- Actually evaluated with statements by the JavaScript interpreter

Examples

• This pattern

- Dynamic with statements
- Dynamically instrumented and executed with statements

Amount

LOADING Set

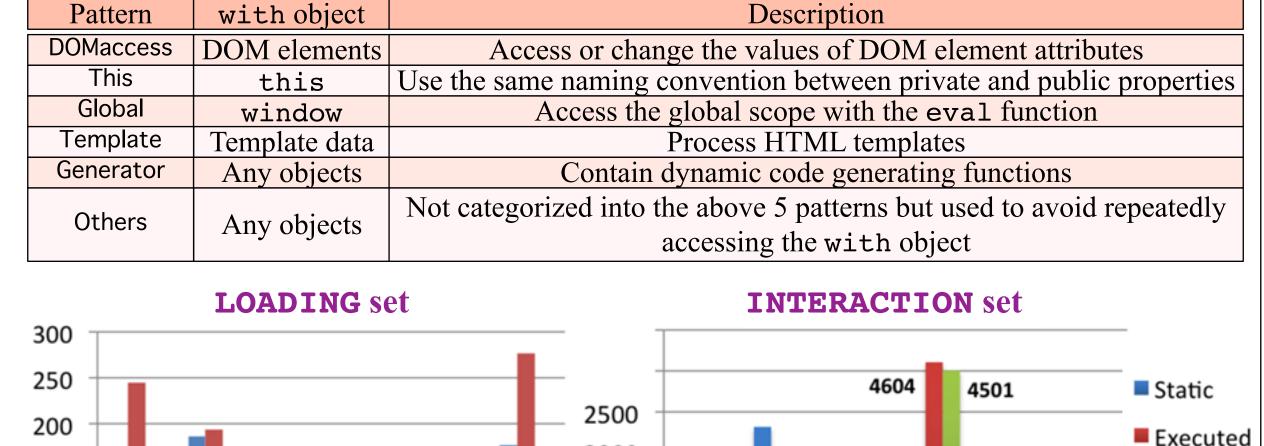
with type	Web sites (944)	with counts	Unique withs
Static	136 (14.4%)	674	645
Executed	66 (6.9%)	855	163
Dynamic	29 (3.0%)	440	75

INTERACTION Set

with type	Web sites (944)	with counts	Unique withs
Static	264 (27.9%)	7,115	1,611
Executed	147 (15.5%)	8,074	1,122
Dynamic	58 (6.1%)	5,665	696

27.9% of the top 944 sites have static occurrences of with statements with simple user interactions. => The with statement is being used unneglectably.

Usage Patterns



with extractor

using the Rhino

JavaScript parser

try { eval(_1f); Executed Dynamic

this.publicF=x; this.publicM=function(){ with(this) privateF=publicF; this.publicF • Global pattern (ebay.com) with(window)

function simpleCons(x){

var privateF;

return true; } catch(e) {}

• The Template pattern (163.com) with(obj){url:"a.com",text:"b" _.push('',text,'');

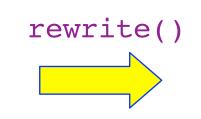
b

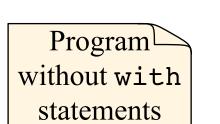
Rewritability

Goal

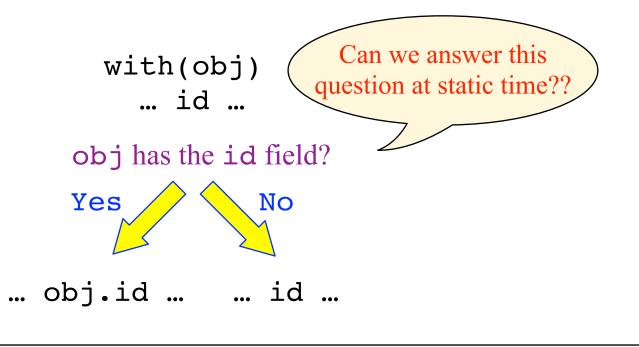
• Defining the **rewrite** function



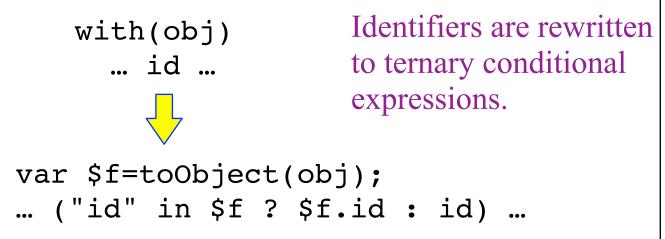




Main Idea



Rewriting the Identifier



Rewriting Rules

150

100

 $rewrite_s[\![with (e) \ s \ | \ \Gamma]\!]$ = If Γ ≡ • Then $\{ \text{var } \alpha = \text{toObject}(rewrite_e[\![e \mid \Gamma]\!]); rewrite_s[\![s \mid \langle \alpha, \bullet, \text{False} \rangle]\!] \}$ Else Let $\Gamma = \langle \phi, \varphi, \beta \rangle$ where $\alpha \notin \phi$ $\{ \text{var } \alpha = \text{toObject}(rewrite_e[\![e \mid \Gamma]\!]); rewrite_s[\![s \mid \langle \phi \mid \alpha, \varphi, \beta \rangle]\!] \}$ $rewrite_{l} \llbracket id \mid \Gamma \rrbracket$ = If $\Gamma \equiv \bullet$ Then idElse Let $\Gamma = \langle \phi \alpha, \varphi, \beta \rangle$ If $id \in \varphi$ Then idElse If $\phi \equiv \bullet$ Then ("id" in α ? α . id: id) Else ("id" in α ? α .id: rewrite_l[[id | $\langle \phi, \varphi, \beta \rangle$]) $rewrite_e \llbracket lhs \odot e \mid \Gamma \rrbracket$ = If $\Gamma \equiv \bullet$ Then $rewrite_{l} \llbracket lhs \mid \Gamma \rrbracket \odot rewrite_{e} \llbracket e \mid \Gamma \rrbracket$ Else Let $\Gamma = \langle \phi \alpha, \varphi, \beta \rangle$ If $lhs \equiv id$ Then If $lhs \in \varphi$ Then $lhs \odot rewrite_e[\![e \mid \Gamma]\!]$ Else If $\phi \equiv \bullet$ Then If $\beta \equiv \text{True}$ Then ("lhs" in α ? α . $lhs \odot e$: $lhs \odot e$) Else ("lhs" in α ? α . lhs \odot rewrite_e [[e | Γ]] : $lhs \odot rewrite_e[\![e \mid \Gamma]\!]$) Else If $\beta \equiv \text{True}$ Then ("lhs" in α ? $\alpha . lhs \odot e$: $rewrite_e[[lhs \odot e \mid \langle \phi, \varphi, \beta \rangle]])$ Else ("lhs" in α ? α . $lhs \odot rewrite_e \llbracket e \mid \Gamma \rrbracket$: $rewrite_e[[lhs \odot rewrite_e[[e \mid \Gamma]]]]$ $\langle \phi, \varphi, \text{True} \rangle]$

Else $rewrite_l \llbracket lhs \mid \Gamma \rrbracket \odot rewrite_e \llbracket e \mid \Gamma \rrbracket$

• Implemented in Java: http://plrg.kaist.ac.kr/research/software

Rewriting the Assignment

Incorrect with(obj) var \$f=toObject(obj); ("x" in \$f ? \$f.x : x)=3;x=3;ReferenceError exception!! Correct rewriting By placing the assignment in

var \$f=toObject(obj); ("x" in f ? f.x=3 : x=3); we can preserve the original

each branch of the conditional. semantics of the with statement.

Rewritability Check

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Pattern	Rewritability	Rewriting
DOMaccess	Yes	By the rewrite function
This	Yes	By the rewrite function
Global	Yes	By ECMAScript 5
Template	Yes	By the rewrite function
Generator	No	Generally not possible due to the dynamic code
		generating functions in the with statement
Others	Yes	By the rewrite function

We can rewrite all static with statements in all patterns except for Generator pattern.

=> 92% of static with statements in the LOADING set and 93% in the INTERACTION set are rewritable.