The Performance Analysis of Garbage Collection on Two Architectures

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1. Introduction

In computer science, garbage collection (GC) is a form of automatic memory management. The garbage collector, referred also to collector, attempts to reclaim garbage or memory occupied by objects that are no longer in use by the program.

Garbage collection was first proposed by John McCarthy around 1959. Garbage collectors for languages like ML and Java, which exert tight control over how applications create and use pointers, can maintain an exact representation of the reachability graph, and thus can reclaim all garbage.

2. Collectors

MarkSweep

During collection

After collection

Reachable (live) object
Unreachable (dead) object
Forwarding object

SemiSpace

From Space

Root Space

To Space

Reference Counting

1. Before decrement(R) on the memory references

2. After decrement(R) on the memory references

3. After decrement(Q) on the memory references, all having count of 0 are reclaimed.

Generational collector

Nursery

Live Objects

Freed objects

Garbage

Old generation

Objects from the previous generation

3. Result

AMD Opteron

Intel core i3

We explores and quantifies garbage collection behavior for three whole heap collectors (MarkSweep, SemiSpace, RefCount) and generational counterparts (GenMS, GenCopy, GenRC). In order to analyze the performance of garbage collectors, we carried out an experiment of garbage collection using DaCapo-9.12-bash benchmark on two architectures which are AMD Opteron and Intel core i3. As a result, it is implied that the how they match program characteristics depending on the resource of hardware and the scale of programs.